

Cryptography and Encryption
Algorithms
CS 239
Computer Security
February 5, 2003

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Outline

- Uses of cryptography
- Symmetric cryptography
- Asymmetric cryptography

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Uses of Cryptography

- What can we use cryptography for?
- Lots of things
 - Secrecy
 - Authentication
 - Prevention of alteration

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Cryptography and Secrecy

- Pretty obvious
- Only those knowing the proper keys can decrypt the message
 - Thus preserving secrecy
- Used cleverly, it can provide other forms of secrecy

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Cryptography and Zero-Knowledge Proofs

- With a really clever use, cryptography can be used to prove I know a secret
 - Without telling you the secret
- Seems like magic, but it can work
- Basically, using multiple levels of cryptography in very clever ways

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Cryptography and Authentication

- How can I prove to you that I created a piece of data?
- What if I give you the data in encrypted form?
 - Using a key only you and I know
- Then only you or I could have created it
 - Unless one of us told someone else the key . . .

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Some Limitations on Cryptography and Authentication

- If both parties cooperative, cryptography can authenticate
 - Problems with non-repudiation, though
- What if three parties want to share a key?
 - No longer certain who created anything
 - Public key cryptography can solve this problem, too
- What if I want to prove authenticity without secrecy?

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Cryptography and Non-Alterability

- Changing one bit of an encrypted message completely garbles it
- If a checksum is part of encrypted data, that's detectable
- If you don't need secrecy, can get the same effect
 - By just encrypting the checksum

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Symmetric and Asymmetric Cryptosystems

- Symmetric - the encrypter and decrypter share a secret key
 - Used for both encrypting and decrypting
- Asymmetric – encrypter has different key than decrypter

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Description of Symmetric Systems

- $C = E(K, P)$
- $P = D(K, C)$
- $E()$ and $D()$ are not necessarily symmetric operations

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Advantages of Symmetric Key Systems

- + Encryption and authentication performed in a single operation
- + Well-known (and trusted) ones perform faster than asymmetric key systems
- + Doesn't require any centralized authority
 - Though key servers help a lot

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Disadvantage of Symmetric Key Systems

- Encryption and authentication performed in a single operation
 - Makes signature more difficult
- Key distribution can be a problem
- Scaling

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Scaling Problems of Symmetric Cryptography

How many keys am I going to need to handle the entire Internet???

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Sample Symmetric Key Ciphers

- The Data Encryption Standard
- The Advanced Encryption Standard
- There are many others

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The Data Encryption Standard

- Probably the best known symmetric key cryptosystem
- Developed in 1977
- Still in wide use
 - Which implies no one has seriously broken it, yet
- But showing its age

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History of DES

- Developed in response to National Bureau of Standards studies
- Developed by IBM
- Analyzed, altered, and approved by the National Security Agency
- Adopted as a federal standard
- One of the most widely used encryption algorithms

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Overview of DES Algorithm

- A block encryption algorithm
 - 64 bit blocks
- Uses substitution and permutation
 - Repeated applications
 - 16 cycles worth
- 64 bit key
 - Only 56 bits really used, though

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More On DES Algorithm

- Uses substitutions to provide confusion
 - To hide the set of characters sent
- Uses transpositions to provide diffusion
 - To spread the effects of one plaintext bit into other bits
- Uses only standard arithmetic and logic functions and table lookup

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Description of DES Algorithm

- Alternate applications of two different ciphers
 - A *product cipher*
- Starts by breaking block in half
- The algorithm goes through 16 *rounds*
- Each round consists of a substitution followed by a permutation

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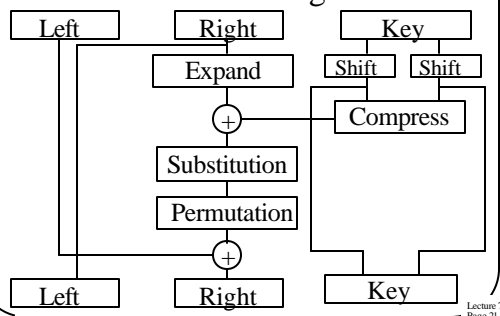
One DES Round

- Select 48 bits from the key
- Expand right half of block to 48 bits
- XOR with key bits
- Look up result in an S-box
 - Resulting in 32 bits
- XOR with left half of block
- Result is new right half
- Old right half becomes new left half

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DES Round Diagram



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S-Boxes

- Table lookups to perform substitutions
- Permanently defined for DES
- Eight different S-boxes
 - Six bits out of 48 bits go to each
 - Four bits come out of each
- Choice of contents of S-boxes believed to strongly impact security of DES

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P-Box

- Maps 32 input bits to 32 output bits
- A single, straight permutation
 - Unlike S-boxes, which are table lookups

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Decrypting DES

- For DES, $D()$ is the same as $E()$
- You decrypt with exactly the same algorithm
- If you feed ciphertext and the same key into DES, the original plaintext pops out

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Is DES Secure?

- Apparently, reasonably
- No evidence NSA put a trapdoor in
- Some keys are known to be weak with DES
 - So good implementations reject them
- Initially, only brute force attacks have cracked DES

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Key Length and DES

- Easiest brute force attack is to try all keys
 - Looking for a meaningful output
- Cost of attack proportional to number of possible keys
- Is 2^{56} enough keys?

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DES Cracking Experiments

- RSA Data Security issued challenge to crack a DES-encrypted message
- Various people got together to do so
 - Harnessing computers across the Internet
 - Using a brute-force approach
- Did it in 1998

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How the DES Message Was Cracked

- Required use of tens of thousands of computers
- Took four months
- The searchers “got lucky”
 - Only one quarter of key space searched
 - On average, brute force requires searching one half of key space
- Done over four years ago
 - So it would presumably take 1/4 time today

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Does This Mean DES is Unsafe?

- Depends on what you use it for
- In how many cases will tens of thousands of machines apply spare cycles for several days to break one message?
- On the other hand, computers will continue to get faster
- And motivated opponents can harness vast resources
- Those who care seriously about security don't tend to use DES any more

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Triple DES

- Simple way of increasing security of DES
- Apply DES three times iteratively to each block
 - Thus, 1/3 as fast as DES
- Use different key for each encryption
- Effectively doubles the key length of DES
- Approved by NIST
 - Which recommends using in preference to DES

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The Advanced Encryption Standard

- A relatively new cryptographic algorithm
- Intended to be the replacement for DES
- Chosen by NIST
 - Through an open competition
- Chosen cipher was originally called Rijndael
 - Developed by Dutch researchers
 - Uses combination of permutation and substitution

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Public Key Encryption Systems

- The encrypter and decrypter have different keys

$$C = E(K_E, P)$$

$$P = D(K_D, C)$$

- Often, works the other way, too

$$C \stackrel{?}{=} E(K_D, P)$$

$$P \stackrel{?}{=} D(K_E, C)$$

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History of Public Key Cryptography

- Invented by Diffie and Hellman in 1976
- Merkle and Hellman developed Knapsack algorithm in 1978
- Rivest-Shamir-Adelman developed RSA in 1978
 - Most popular public key algorithm

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Practical Use of Public Key Cryptography

- Keys are created in pairs
- One key is kept secret by the owner
- The other is made public to the world
- If you want to send an encrypted message to someone, encrypt with his public key
 - Only he has private key to decrypt

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Authentication With Shared Keys

- If only two people know the key, and I didn't create a properly encrypted message -
 - The other guy must have
- But what if he claims he didn't?
- Or what if there are more than two?
- Requires authentication servers

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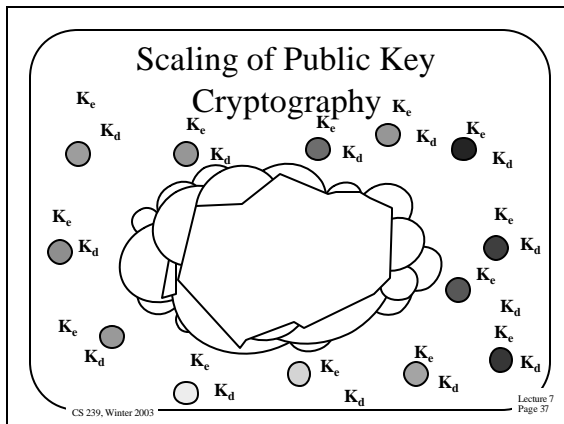
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Authentication With Public Keys

- If I want to "sign" a message, encrypt it with my private key
- Only I know private key, so no one else could create that message
- Everyone knows my public key, so everyone can check my claim directly

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Key Management Issues

- To communicate via shared key cryptography, key must be distributed
 - In trusted fashion
- To communicate via public key cryptography, need to find out each other's public key
 - “Simply publish public keys”

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Issues of Key Publication

- Security of public key cryptography depends on using the right public key
- If I am fooled into using the wrong one, that key's owner reads my message
- Need high assurance that a given key belongs to a particular person
- Which requires key distribution infrastructure

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RSA Algorithm

- Most popular public key cryptographic algorithm
- In wide use
- Has withstood much cryptanalysis
- Based on hard problem of factoring large numbers

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RSA Keys

- Keys are functions of a pair of 100-200 digit prime numbers
- Relationship between public and private key is complex
- Recovering plaintext without private key (even knowing public key) is supposedly equivalent to factoring product of the prime numbers

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Comparison of DES and RSA

- DES is much more complex
- However, DES uses only simple arithmetic, logic, and table lookup
- RSA uses exponentiation to very large powers
 - Computationally 1000 times more expensive in hardware
 - 100 times in software
- Key selection also more expensive

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A Downside To RSA

- It's patented
- By the inventor's company
- Thus, you need a license to use it
- Company has had some criticism on licensing policies
- But very widely used today

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Security of RSA

- Conjectured that security depends on factoring large numbers
 - But never proven
 - Some variants proven equivalent to factoring problem
- Probably the conjecture is correct

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Attacks on Factoring RSA Keys

- In 1994, a 129 bit RSA key was successfully factored
 - 1600 computers, several months
- Recent paper (fall 2001) on integer factorization suggests keys up to 2048 bits may be insecure
- Size will keep increasing
- The longer the key, the more expensive the encryption and decryption

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Combined Use of Symmetric and Asymmetric Cryptography

- Very common to use both in a single session
- Asymmetric cryptography essentially used to “bootstrap” symmetric crypto
- Use RSA (or other algorithm like PGP) to authenticate and establish a session key
- Use DES/Triple DES/AES using session key for the rest of the transmission

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Digital Signature Algorithms

- In some cases, secrecy isn't required
- But authentication is
- The data must be guaranteed to be that which was originally sent
- Especially important for data that is long-lived

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Desirable Properties of Digital Signatures

- Unforgeable
- Verifiable
- Non-repudiable
- Cheap to compute and verify
- Non-reusable
- No reliance on trusted authority
- Signed document is unchangeable

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Encryption and Digital Signatures

- Digital signature methods are based on encryption
- Encryption can be used as a signature

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Signatures With Shared Key Encryption

- Requires a trusted third party
- Signer encrypts document with secret key shared with third party
- Receiver checks validity of signature by consulting with trusted third party
- Third party required so receiver can't forge the signature

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Signatures With Public Key Cryptography

- Signer encrypts document with his private key
- Receiver checks validity by decrypting with signer's public key
- Only signer has the private key
 - So no trusted third party required
- But receiver must be certain that he has the right public key

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Problems With Simple Encryption Approach

- Computationally expensive
 - Especially with public key approach
- Document is encrypted
 - Must be decrypted for use
 - If in regular use, must store encrypted and decrypted versions

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Secure Hash Algorithms

- A method of protecting data from modification
- Doesn't actually prevent modification
- But gives strong evidence that modification did or didn't occur
- Typically used with digital signatures

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Idea Behind Secure Hashes

- Apply a one-way cryptographic function to data in question
- Producing a much shorter result
- Attach the cryptographic hash to the data before sending
- When necessary, repeat the function on the data and compare to the hash value

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Secure Hash Algorithm (SHA)

- Endorsed by NIST
- But produced by the NSA . . .
- Reduces input data of up to 2^{64} bits to 160 bit digest
- Doesn't require secret key
- Generally felt to be reasonably secure

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Use of Cryptographic Hashes

- Must assume opponent also has hashing function
- And it doesn't use secret key
- So opponent can substitute a different message with a different hash
- How to prevent this?
- And what (if anything) would secure hashes actually be useful for?

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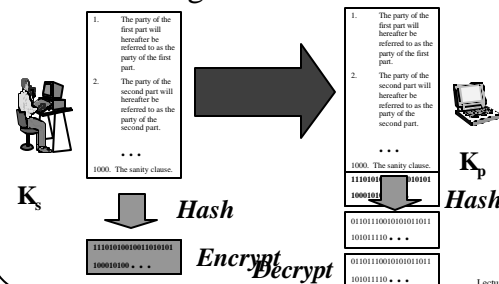
Hashing and Signatures

- Use a digital signature algorithm to sign the hash
- But why not just sign the whole message, instead?
- Computing the hash and signing it may be faster than signing the document
- Receiver need only store document plus hash

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Checking a Document With a Signed Hash



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The Birthday Attack

- How many people must be in a room for the chances to be greater than even that two of them share a birthday?
- Answer is 23
- The same principle can be used to attack hash algorithms

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Using the Birthday Attack on Hashes

- For a given document, find a different document that has the effect you want
- Trivially alter the second document so that it hashes to the same value as the target document
 - Using an exhaustive attack

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How Hard Is the Birthday Attack?

- Depends on the length of the hash
 - And the quality of the hashing algorithm
- Essentially, looking for hashing collisions
- So long hashes are good
 - SHA produces 2^{80} random hashes