

# Authentication

## CS 239

### Computer Security

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## Authentication for Single Machines

- Most single machine system security mechanisms are based on controlling access
- Access control only works if you have good authentication
- Various means are used to provide authentication in operating systems

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## Process Authentication

- Memory protection is based on process identity
  - Only the owning process can name its own virtual memory pages
- Because VM is completely in OS control, pretty easy to ensure that processes can't fake identities

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## How the OS Authenticates Processes

- System calls are issued by a particular process
- The OS securely ties a process control block to the process
  - Not under user control
- Thus, the ID in the process control block can be trusted

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## How Do Processes Originally Obtain Access Permission?

- Most OS resources need access control based on user identity or role
  - Other than virtual memory pages and other transient resources
- How does a process get properly tagged with its owning user or role?
- Security is worthless if OS carefully controls access on a bogus user ID

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## Users and Roles

- In most systems, OS assigns each potential user an ID
- More sophisticated systems recognize that the same user works in different *roles*
  - Effectively, each role requires its own ID
  - And secure methods of setting roles

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## Securely Identifying Users and Roles

- Passwords
- Identification devices
- Challenge/response systems
- Physical verification of the user

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## Passwords

- Authentication by what you know
- One of the oldest and most commonly used security mechanisms
- Authenticate the user by requiring him to produce a secret
  - Known only to him and to the authenticator
  - Or, if one-way encryption used, known only to him

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## Problems With Passwords

- They have to be unguessable
  - Yet easy for people to remember
- If networks connect terminals to computers, susceptible to password sniffers
- Unless fairly long, brute force attacks often work on them

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## Proper Use of Passwords

- Passwords should be sufficiently long
- Passwords should contain non-alphabetic characters
- Passwords should be unguessable
- Passwords should be changed often
- Passwords should never be written down
- Passwords should never be shared

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## Passwords and Single Sign-On

- Many systems ask for password once
  - Resulting authentication lasts for an entire “session”
- Unless other mechanisms in place, complete mediation definitely not achieved
- Trading security for convenience

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## Handling Passwords

- The OS must be able to check passwords when users log in
- So must the OS store passwords?
- Not really
  - It can store an encrypted version
- Encrypt the offered password
  - Using a one-way function
- And compare it to the stored version

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### Standard Password Handling

The Marx Brothers' Family Machine

Harpo	2st6'sG0
Zeppo	G>15{as3
Chico	w*:-sddw
Karl	sY(34,ee
Groucho	We6/d02,
Gummo	3(wbnP1

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### Is Encrypting the Password File Enough?

- What if an attacker gets a copy of your password file?
- No problem, the passwords are encrypted
  - Right?
- Yes, but . . .

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### Dictionary Attacks on an Encrypted Password File

Harpo	2st6'sG0
Zeppo	G>15{as3
Chico	sY(34,ee
Karl	3(wbnP1
Groucho	sY(34,ee
Gummo	3(wbnP1

Now you can hack the Communist Manifesto! **Rats!!!!**

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### A Serious Issue

- All Linux machines use the same one-way function to encrypt passwords
- If someone runs the entire dictionary through that function,
  - Will they have a complete list of all encrypted dictionary passwords?

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### Illustrating the Problem

Darwin: Λ\*eP61a-  
 Marx: Λ\*eP61a-  
 Dictionary: aardvark, aardwolf, 340ja, K[ds, sY(34,ee, beard, Λ\*eP61a-

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### The Real Problem

- Not that Darwin and Marx chose the same password
- But that anyone who chose that password got the same encrypted result
- So the attacker need only encrypt every possible password once
- And then she has a complete dictionary usable against anyone

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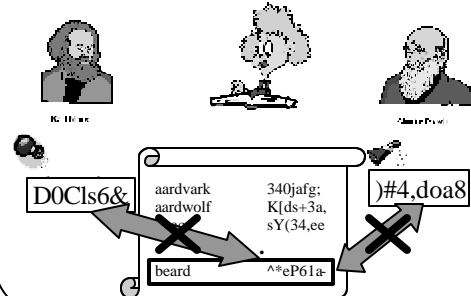
## Salted Passwords

- Combine the plaintext password with a random number
  - Then run it through the one-way function
- The random number need not be secret
- It just has to be different for different users

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## Did It Fix Our Problem?



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## Protecting the Password File

- So it's OK to leave the encrypted version of the password file around?
- No, it isn't
- Why make it easy for attackers?
- Dictionary attacks against single accounts can still work
- Generally, don't give access to the encrypted file, either

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## Identification Devices

- Authentication by what you have
- A smart card or other hardware device that is readable by the computer
- Authenticate by providing the device to the computer

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## Problems With Identification Devices

- If lost or stolen, you can't authenticate yourself
  - And someone else can
  - Often combined with passwords to avoid this problem
- Unless cleverly done, susceptible to sniffing attacks
- Requires special hardware

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## Challenge/Response Authentication

- Authentication by what questions you can answer correctly
- The system asks the user to provide some information
- If it's provided correctly, the user is authenticated

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### Differences From Passwords

- Challenge/response systems ask for different information every time
- Or at least the questions come from a large set
- Best security achieved by requiring what amounts to encryption of the challenge
  - But that requires special hardware
  - Essentially, a smart card

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### Problems With Authentication Through Challenge/Response

- Either the question is too hard to answer without special hardware
- Or the question is too easy for intruders to spoof the answer
- Still, commonly used in real-world situations
  - E.g., authenticating you by asking your mother's maiden name

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### Authentication Through Physical Verification

- Authentication based on who you are
- Things like fingerprints, voice patterns, retinal patterns, etc.
- To authenticate to the system, let it measure the appropriate physical characteristics

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### Problems With Physical Verification

- Requires very special hardware
  - Possibly excepting systems that examine typing patterns
- May not be as foolproof as you think
- Many characteristics vary too much for practical use
- Generally not helpful for authenticating programs or roles
- What happens when it's cracked?
  - You only have two retinas, after all

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### Authenticating Across the Network

- What new challenges does this add?
- You don't know what's at the other end of the wire
- So, when does that cause a problem?
- And how can you solve it?

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