

# Security Principles and Mechanisms

## CS 239

### Computer Security

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## Outline

- Security terms and concepts
- Mechanisms
- Access control

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## Security and Protection

- *Security* is a policy
  - E.g., “no unauthorized user may access this file”
- *Protection* is a mechanism
  - E.g., “the system checks user identity against access permissions”
- Protection mechanisms implement security policies

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## Design Principles for Secure Systems

- Economy
- Complete mediation
- Open design
- Separation of privileges
- Least privilege
- Least common mechanism
- Acceptability
- Fail-safe defaults

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## Economy in Security Design

- Economical to develop
  - And to use
  - And to verify
- Should add little or no overhead
- Should do only what needs to be done
- Generally, try to keep it simple and small

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## Complete Mediation

- Apply security on every access to a protected object
  - E.g., each read of a file, not just the open
- Also involves checking access on everything that could be attacked

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## Open Design

- Don't rely on "security through obscurity"
- Assume all potential attackers know everything about the design
  - And completely understand it
- This doesn't mean publish everything important about your security system
  - Though sometimes that's a good idea

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## Separation of Privileges

- Provide mechanisms that separate the privileges used for one purpose from those used for another
- To allow flexibility in security systems
- E.g., separate access control on each file

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## Least Privilege

- Give bare minimum access rights required to complete a task
- Require another request to perform another type of access
- E.g., don't give write permission to a file if the program only asked for read

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## Least Common Mechanism

- Avoid sharing parts of the security mechanism
  - among different users
  - among different parts of the system
- Coupling leads to possibilities security breaches

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## Acceptability

- Mechanism must be simple to use
- Simple enough that people will use it without thinking about it
- Must rarely or never prevent permissible accesses

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## Fail-Safe Designs

- Default to lack of access
- So if something goes wrong or is forgotten or isn't done, no security lost
- If important mistakes are made, you'll find out about them
  - Without loss of security
  - But if it happens too often . . .

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## Tools for Security

- Physical security
- Access control
- Encryption
- Authentication
- Encapsulation
- Intrusion detection
- Common sense

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## Physical Security

- Lock up your computer
  - Actually, sometimes a good answer
- But what about networking?
  - Networks poke a hole in the locked door
- In any case, lack of physical security often makes other measures pointless

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## Access Controls

- Only let authorized parties access the system
- A lot trickier than it sounds
- Particularly in a network environment
- Once data is outside your system, how can you continue to control it?
  - Again, of concern in network environments

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## Encryption

- Algorithms to hide the content of data or communications
- Only those knowing a secret can decrypt the protection
- One of the most important tools in computer security

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## Encryption is Not a Panacea

- Encryption is usually breakable
  - Given enough time and resources
- Encryption can't protect everything
- Encryption is only as good as the security measures that use it

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## Authentication

- Methods of ensuring that someone is who they say they are
- Vital for access control
- But also vital for many other purposes
- Often (but not always) based on encryption

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## Encapsulation

- Methods of allowing outsiders limited access to your resources
- Let them use or access some things
  - But not everything
- Simple, in concept
- Extremely challenging, in practice

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## Intrusion Detection

- All security methods sometimes fail
- When they do, notice that something is wrong
- And take steps to correct the problem
- Reactive, not preventative
  - But unrealistic to believe any prevention is certain
- Must be automatic to be really useful

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## Common Sense

- A lot of problems arise because people don't like to think
- The best security tools generally fail if people use them badly
- If the easiest way in is to fool people, that's what attackers will do

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## The Depressing Truth

- Ultimately, computer security is a losing battle
- Nothing will ever work 100%
- Nothing will work forever
- All your efforts will eventually be undone
- It's like housework – doing it doesn't make the house clean tomorrow, but not doing it guarantees the house is dirty today

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## Access Control

- Security could be easy
  - If we didn't want anyone to get access to anything
- The trick is giving access to only the right people
- How do we ensure that a given resource can only be accessed by the proper people?

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## Goals for Access Control

- Complete mediation
- Least privilege
- Useful in a networked environment
- Scalability
- Cost and usability

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## Access Control Mechanisms

- Directories
- Access control lists
- Capabilities
- Access control matrices

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## Directories

- Each user has a list of the items he can access
  - With the associated rights
- When a user wants to access an item, look it up in his directory

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## Problems With the Directory Approach

- Per-user directories get very large
  - Overhead and performance problems
- Universal revocation of access
- Pseudonym problems
- Works poorly in networks
- This method is not widely used

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## Access Control Lists

- For each protectable resource, maintain a single list
- Each list entry specifies a user who can access the resource
  - And the allowable modes of access
- When a user requests access to a resource, check the access control list (ACL)

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## ACL Objects and Subjects

- In ACL terminology, the resources being protected are *objects*
- The entities attempting to access them are *subjects*
  - Allowing finer granularity of control than per-user

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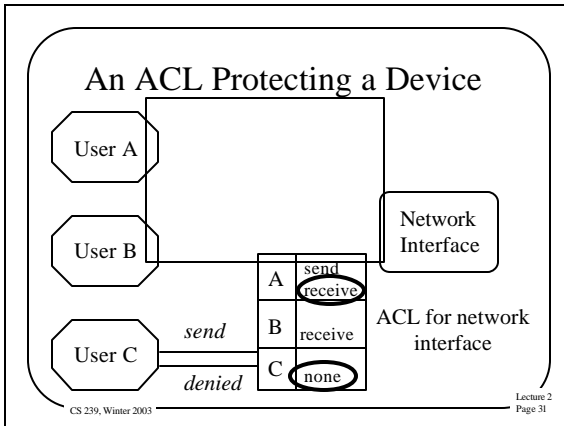
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## ACL Example

- An operating system example:
  - Using ACLs to protect a network interface device
- User A is allowed to receive from and send to the device
- User B may only receive from it
- User C may not access it

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### Issues for Access Control Lists

- How do you know the requestor is who he says he is?
- How do you protect the access control list from modification?
- How do you determine what resources a user can access?

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### ACLs in Practice

- Unix file permissions are a limited form of an ACL
  - Only owner, group, and all can have ACL entries
  - Only read/write/execute controls are available
- Other systems (like Windows NT) have more general ACL mechanisms

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### Pros and Cons of ACLs

- + Easy to figure out who can access a resource
- + Easy to revoke or change access permissions
  - Hard to figure out what a subject can access
  - Changing access rights requires getting to the object

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### Capabilities

- Each subject keeps a set of data items that specify his allowable accesses
- Essentially, a set of tickets
- Possession of the capability for an object implies that access is allowed

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### Properties of Capabilities

- Must be unforgeable
  - In single machine, keep under control of OS
  - What about in a networked system?
- In most systems, some capabilities allow creation of other capabilities
  - Process can pass restricted set of capabilities to a subprocess

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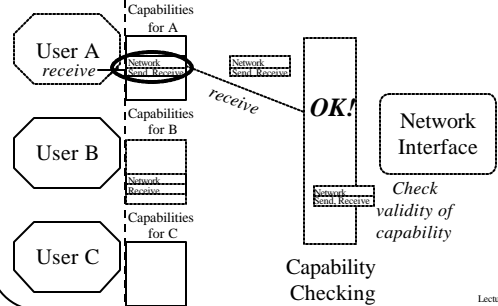
## Capabilities and Domains

- The set of objects a subject can access at a given moment is its domain
  - The subject has a capability for each object in its domain
- Domains can be expanded by obtaining new capabilities
- New domains can be created for subprocesses
- Where do we keep capabilities?

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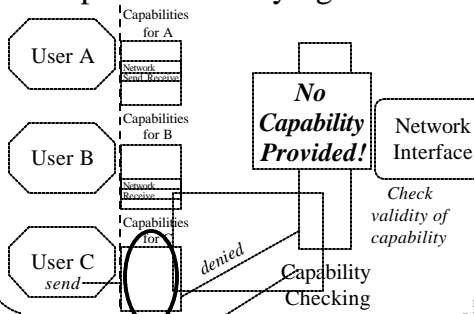
## Capabilities Protecting a Device



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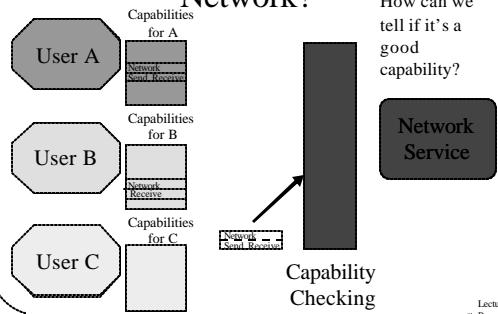
## Capabilities Denying Access



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## How Will This Work in a Network?



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## Revoking Capabilities

- A major challenge in capability systems
- Several methods available:
  - 1). Search and destroy
  - 2). Invalidation at use
  - 3). Indirection through a token
  - 4). Generation numbers

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## Revocation By Destroying Capabilities

- Find the capability you want to revoke
- Destroy it
- Easy if all capabilities live only in system-controlled memory
- But most systems allow storage of capabilities on disk
- And what about networked systems?

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### Revocation By Invalidation on Use

- Keep a list of revoked capabilities
  - Usually one list per object
- When a capability is presented for use, check it against the list
- Expensive, especially if the list is long or complete mediation is used
  - Not feasible on every access
- And what about networked systems?

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### Revocation by Indirection Through a Token

- Capability points to token under system control
- Token is set up on first access to object
- To revoke access, destroy the token
- Adds cost to checking access
- Usually hard to provide selective revocation

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### Revocation By Generation Numbers

- Store a random number in each capability
- Store the same random number with the protected object
- On access, check the numbers
- To revoke access, change the number
- No selective revocation
- Requires some control of capabilities

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### Pros and Cons of Capabilities

- + Easy to determine what a subject can access
- + Potentially faster than ACLs (in some circumstances)
- + Easy model for transfer of privileges
  - Hard to determine who can access an object
  - Requires extra mechanism to allow revocation
  - In network environment, need cryptographic methods to prevent forgery

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### ACLs, Capabilities, Complete Mediation, & Performance

- Ideally, every data access should have access control independently applied
- Practicality of doing so depends on the performance costs
- What does it cost to use ACLs?
  - Capabilities?

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### Performance Issues of Access Control

- What if the status of the access control mechanism changed between when last checked and current access?
- Common case is nothing changes
- Different approaches possible
  - Actually check changeable data structure on each access
  - Give process something cheap and revocable that allows access

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## Access Control and ACLs

- The ACL is a list
- Initially, checking an ACL involves searching a list
- For later checks, maintain pointer to list entry
- Be sure that changing the permissions changes what's pointed to

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## Access Controls and Capabilities

- Attach the capability (or pointer to it) to each request
- Use attached information to determine if current access is permissible
- This approach is hard to use with revocation

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## An Alternate Approach To Using Capabilities

- On first access, use a capability to obtain an access token
  - Using careful, expensive checks to see if capability was revoked
- If revocation required, destroy the access token
- Can also be done with pointers

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## Access Control in the Distributed World

- ACLs still work OK
  - Provided you have a global namespace for subjects
- Capabilities are more problematic
  - Their security relies on unforgeability

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## Using Cryptographic Capabilities

- Can cryptography make capabilities unforgeable?
- It can make it impossible to create them from nothing
  - And only usable by their owner
- But it can't make them uncopyable
- So cryptographic capability systems must assume they can be freely copied

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## Access Control Matrices

- A very general access control concept
- In principle, ACLs are a 1-D list of who is permitted to access one object
- And capabilities are a 1-D list of what one subject can access
- Access control matrices are a 2-D description of access rights

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### Access Control Matrix Example

	File A	File B	Network	Printer	<b>Objects</b>
User 1	rw	r		w	<b>User 2's Capabilities</b>
User 2	r		sr	w	
Sysadmin	rw	rw	sr	rw configure	
Guest			sr		
<b>Subjects</b>	<b>File B's ACL</b>				

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### Pros and Cons of Access Control Matrices

- + Makes all access issues explicit and easy to find
- + Easy to tell who can access a resource, and what resources anyone can access
- Matrix very sparse, so inefficient
- Hard to achieve good performance
- More important conceptually than in implementations

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### Role Based Access Control

- Not really an alternative to ACLs, capabilities, access control matrix
- Rather, a more complex way of looking at access control subjects
- Commonly used in systems that care about security

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### The Idea Behind Role Based Access Control

- Each user has certain roles he can take while using the system
- At any given time, the user is performing a certain role
- Give the user access to only those things that are required to fulfill that role

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### A Simple Example

- Fred is a system administrator
  - Which requires him to install programs, examine logs, etc.
- Fred also reads email, looks at web sites, etc.
- Fred should operate under one role while doing normal work
  - And a different role while performing administrative tasks

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### Continuing With the Example

- Fred logs on as “fred”
- He reads his email as “fred”
- He decides to upgrade the C++ compiler
  - So he changes roles to “administrator”
- When he’s done, he returns to the role of “fred”

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### What Has Been Gained?

- While reading mail and surfing the web, Fred isn't able to upgrade the C++ compiler
  - He doesn't have the access rights
- So if he accidentally downloads malicious code, it can't "upgrade" the compiler

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### Changing Roles

- Role based access control only helps if changing roles isn't trivial
  - Otherwise, the malicious code merely changes roles before doing anything else
- Typically requires providing some secure form of authentication
  - Which proves you have the right to change roles

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### Practical Limitations on Role Based Access Control

- Number of roles per user
- Problems of disjoint role privileges
- System administration overheads

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### Number of Roles Per User

- Each new role requires new authentication
- Less secure if the authentication is the same for each role
  - E.g., Unix `sudo`, which only requires your basic password
- How many passwords will people remember?
  - And how often will they be happy to type them?

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### Problems of Disjoint Roles

- Each role should have disjoint privileges
  - More secure if roles aren't supersets of other roles
- May cause difficulties if certain operations require privileges from different roles

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### Problems of System Administration

- Access control is only useful if the permissions are set correctly for each subject and object
- The more subjects there are, the more work system administrators must do
  - Since each subject needs to get only the proper privileges

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## Discretionary Access Control

- Individual subjects are permitted to decide on access control issues
- And can change them whenever they please
  - Though only for objects they own or control

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## Mandatory Access Control

- A system-wide policy on access control is enforced
- Subjects are not necessarily allowed to alter access controls
  - Even on their own stuff
- Important for organizations that care strongly about security

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## Conclusion

- Much of security relates to allowing some people access to some resources
- While preventing the same access to others
- Without some method of determining who should access what . . .  
You can't do that

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