

Security Alert Systems and
Handling Compromised Hosts
CS 239
Advanced Topics in Network
Security
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Today's Topics

- Bad things happen on the net
 - Many affect almost everyone
 - How do we get the word out?
- Many machines attached to the net are compromised
 - Attackers can do arbitrary bad things with them
 - How do we deal with that situation?

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Security Alert Systems

- Most network problems threaten lots of people
 - Thousands to millions
- Often there are things one can do to protect oneself
- How do we tell people what they should do?

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Difficult Issues for Alert Systems

- Scale
- Delay
- Security
- Trust
- Cost
- Deployment

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Scale and the Alert Problem

- Many threats against Windows
- Tens of millions of Windows machines
- Increasingly, usually attached to the network
- Security alerts must work at that scale

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Delay and the Alert Problem

- Worms can spread in a few minutes
- Unprotected machines attached to the net are compromised in minutes
 - Sometimes seconds
- Shrinking delays between flaws being revealed and being exploited
- How fast do you need to get alerts out?

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Security and the Alert Problem

- Alert systems make recipients do something
- Actions can have wide-ranging effects
 - Close a port?
 - Refuse some types of email?
 - Refuse to talk to certain machines?
- What if attackers generate an alert?
- Can they make large numbers of nodes behave badly?
- How do we stop that from happening?

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Trust and Alert Systems

- Assuming we have security,
- Only trusted parties send these powerful alerts
- But who should we trust?
- The “obvious suspects” have done things that suggest trusting them might not be good
 - Not alerting to all problems
 - Not sending alerts very quickly
 - Actions taken for their own motives

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Cost and Alert Systems

- Mostly meant to protect us from hazards
- Hazards don't come up that often
- How much should we pay to get alerts?
 - At alert time?
 - At all times?
 - And where?
- If there's a monetary cost, who pays?
 - And why would they pay?

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Deployment and Alert Systems

- Where do alert components need to be deployed?
- Why would owners of deployment points cooperate?
- How well does the system work if they don't?
 - Never?
 - At critical moments?

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Basic Approaches to Security Alerts

- Centralized alert servers
 - CERT, Symantec, Microsoft, etc.
- Distributed alert servers
 - A la DNS servers
- Peer systems
 - Revere
 - DHT-based
 - Unstructured

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What Gets Alerted?

- Worm spread
- Malware signatures
- Sharing of spam signatures
- Blacklists of sites
- Patches
 - SW and firewall rules
- What about use for non-cyber alerts?
 - Tornados, hurricanes, tsunami?
- One alert system for everything?
 - Or several different systems?

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A Different Approach

- Don't rely on SW infrastructure to spread alerts
- Rely on people
 - Send out to mailing lists
 - Post on web sites
 - Assume people will tell their friends
- How is this better (or worse)?
- How can we “bundle” this effect with pure SW distribution systems?

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Handling Compromised Machines

- Many machines are compromised
 - Hundreds of thousands at least
 - Maybe millions
- Many of them are never fixed
- Many of them are permanently attached to Internet
- What problems does this cause?
- What can we do about them?

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Typical Actions of Compromised Machines

- “Manually” compromising other machines
- Spam agents
- Phishing agents
- Storing copyrighted material (warez)
- DDoS agents
- Harvesting of sensitive material
- Password cracking
- Other unforeseen uses

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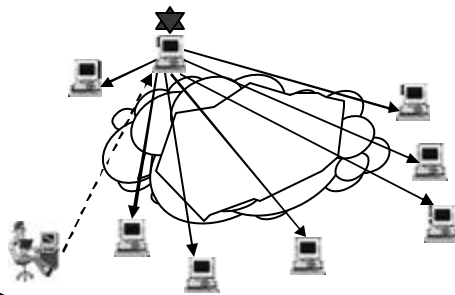
Botnets

- A recent development
 - Three or four years old, mostly
- Attempts by attackers to organize sets of compromised machines
 - Using distributed system techniques
- Eases management
- Allows automated control of large numbers of machines

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Typical Botnet Organization



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Specifics of Typical Botnet Organization

- An overlay network
 - Generally not optimized for underlying network
- Usually hierarchically arranged
 - Generally not deep hierarchy
- Often little resiliency
 - Bots know about 1 master
 - But usually mechanisms to change that

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Typical Botnet Organization

- Communication generally over IRC
- Usually master/slave communications
 - Usually one master machine
 - Master usually identified by password/key
 - Not by node identity
 - Big botnets might require hierarchy

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Approaches to Handling Compromised Machines

- Force cleanup
 - How?
 - “Attack-back” or forced patching legally questionable
- Identify and blacklist compromised nodes
 - Or treat them with caution
- Live with it

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Blacklist Approaches

- Someone makes list of compromised nodes
- Notifies other participants in system
- All participants refuse traffic from blacklisted sites
 - Or treat it cautiously

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Challenges for Blacklists

- How do you identify compromised nodes?
- How do you alert others?
- What if you’re wrong?
- How do cleaned-up nodes get off the blacklist?
- Boils down to some typical distributed systems/data replication questions

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Handling Botnets

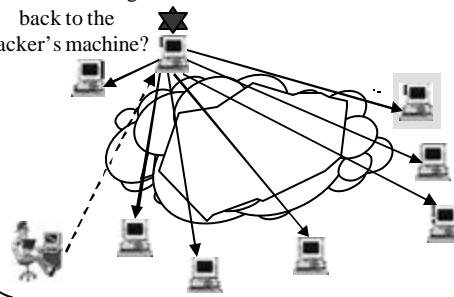
- How do you figure out which nodes are in a botnet?
- How do you interfere with a botnet’s operations?
- How do you trace from botnet to attacker’s machine?

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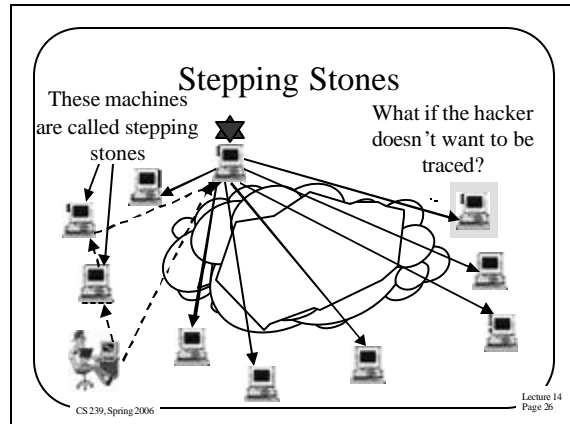
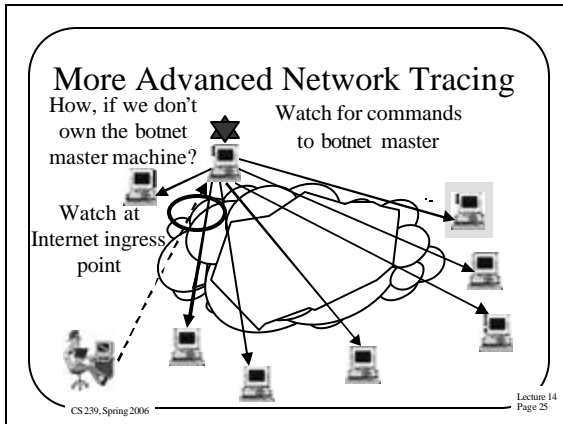
Basic Network Tracing

But how do we get
back to the
hacker’s machine?



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- ### Complexities of Dealing With Stepping Stones
- Requires tracing at multiple points
 - Can you get to all of them?
 - Complicates tracing by confounding traffic
 - May be hard to pull command traffic out of all stepping stone traces
 - Makes it hard to figure out how far you have to go
 - Takes longer
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