

Network Security
CS 239
Security for Networks and
System Software
April 29, 2002

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Outline

- Catching up on certificates
- Basics of network security

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Certificates

- An increasingly popular form of authentication
- Generally used with public key cryptography
- A signed electronic document proving you are who you claim to be

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Public Key Certificates

- The most common kind of certificate
- Addresses the biggest challenge in widespread use of public keys
- Essentially, a copy of your public key signed by a trusted authority
- Presentation of the certificate alone serves as authentication of your public key

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Implementation of Public Key
Certificates

- Set up a universally trusted authority
- Every user presents his public key to the authority
- The authority returns a certificate
 - Containing the user's public key signed by the authority's private key

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Checking a Certificate

- Every user keeps a copy of the authority's public key
- When a new user wants to talk to you, he gives you his certificate
- Decrypt the certificate using the authority's public key
- You now have an authenticated public key for the new user
- Authority need not be checked on-line

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Scaling Issues of Certificates

- If there are ~550 million Internet users needing certificates, can one authority serve them all?
- Probably not
- So you need multiple authorities
- Does that mean everyone needs to store the public keys of all authorities?

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Certification Hierarchies

- Arrange certification authorities hierarchically
- The single authority at the top produces certificates for the next layer down
- And so on, recursively

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Using Certificates From Hierarchies

- I get a new certificate
- I don't know the signing authority
- But the certificate also contains that authority's certificate
- Perhaps I know the authority who signed this authority's certificate

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Extracting the Authentication

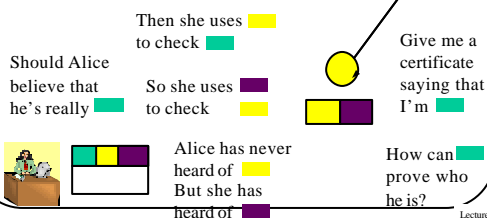
- Using the public key of the higher level authority, extract the public key of the signing authority from the certificate
- Now I know his public key, and it's authenticated
- I can now extract the user's key and authenticate it

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A Example

Alice gets a message with a certificate



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Certificates and Trust

- Ultimately, the point of a certificate is to determine if something is trusted
 - Do I trust the request to perform some financial transaction?
- So, Trustysign.com signed this certificate
- How much confidence should I have in the certificate?

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Potential Problems in the Certification Process

- What measures did Trustysign.com use before issuing the certificate?
- Is the certificate itself still valid?
- Is Trustysign.com's signature/certificate still valid?
- Who is trustworthy enough to be at the top of the hierarchy?

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Trustworthiness of Certificate Authority

- How did Trustysign.com issue the certificate?
- Did it get an in-person sworn affidavit from the certificate's owner?
- Did it phone up the owner to verify it was him?
- Did it just accept the word of the requestor that he was who he claimed to be?

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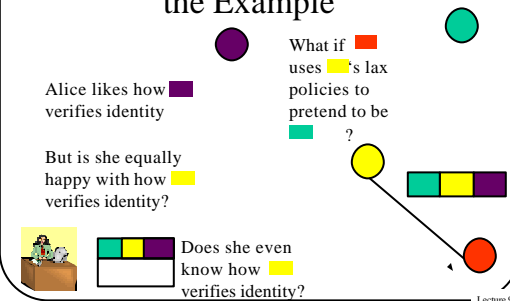
What Does a Certificate Really Tell Me?

- That the certificate authority (CA) tied a public/private key pair to identification information
- Generally doesn't tell me why the CA thought the binding was proper
- I may have different standards than that CA

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Showing a Problem Using the Example



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Another Big Problem

- Things change
- One result of change is that what used to be safe or trusted isn't any more
- If there is trust-related information out in the network, what will happen when things change?

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Revocation


- A general problem for keys, certificates, access control lists, etc.
- How does the system revoke something related to trust?
- In a network environment
- Safely, efficiently, etc.

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Revisiting Our Example

Someone discovers that  has obtained a false certificate for 

How does Alice make sure that she's not accepting 's false certificate?



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The Web of Trust Model

- Public keys are still passed around signed by others
- But your trust in others is based on your personal trust of them
 - Not on a formal certification hierarchy
 - “I work in the office next to Bob, so I trust Bob’s certifications”

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Certificates in the Web of Trust

- Any user can sign any other user’s public key
- When a new user presents me his public key, he gives me one or more certificates signed by others
- If I trust any of those others, I trust the new user’s public key

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Limitations on the Web of Trust

- The web tends to grow
 - “I trust Alice, who trusts Bob, who trusts Carol, who trusts Dave, . . . , who trusts Lisa, who trusts Mallory”
 - Just because Lisa trusts Mallory doesn’t mean I should
- Working system needs concept of degrees of trust

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Advantages and Disadvantages of Web of Trust Model

- + Scales very well
- + No central authority
- + Very flexible
 - May be hard to assign degrees of trust
 - Revocation may be difficult
 - May be hard to tell who you will and won’t trust

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Some Important Network Characteristics for Security

- Degree of locality
- Media used
- Protocols used

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Degree of Locality

- Some networks are very local
 - E.g., an Ethernet
 - Only handles a small number of machines, mostly related ones
- Other networks are very non-local
 - E.g., the Internet backbone
 - Vast numbers of users/sites share bandwidth

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Implications of Locality

- Truly local networks may gain from physical security
- Relative trustworthiness of all participants may help
- Common interests of all on a local network may be helpful, too
- Wide area networks generally harder

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Network Media

- Some networks are wires or cables
- Other networks run over the telephone lines
- Other networks are radio links to satellites
- Other networks are broadcast radio links

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Implications of Media Type

- Wires can sometimes be physically protected
- Radio links generally can't
 - Though power and technology requirements for satellite links may provide some help

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Protocol Types

- TCP/IP is probably the most widespread
 - But it only specifies some common intermediate levels
 - Other protocols exist above and below it
- And, in places, other protocols replace TCP/IP
- And there are lots of supporting protocols
 - Routing protocols, naming and directory protocols, network management protocols

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Implications of Protocol Type

- The protocol defines a set of rules that will always be followed
 - But usually not quite complete
 - And they assume everyone is at least trying to play by the rules
 - What if they don't?
- Specific attacks exist against specific protocols

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Threats to Network Security

- Pretty much the usual suspects:
 - Wiretapping
 - Impersonation
 - Message confidentiality
 - Message integrity
 - Denial of service

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Why Are Networks Especially Threatened?

- Many “moving parts”
- Many different administrative domains
- Everyone can get some access
- In some cases, trivial for attacker to get a foothold on the network
- Networks encourage sharing
- Networks often allow anonymity

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What Can Attackers Attack?

- The media connecting the nodes
- Nodes that are connected to them
- Routers that control the traffic
- The protocols that set the rules for communications

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Wiretapping

- An obvious network vulnerability
 - But don’t forget, “wiretapping” is a general term
 - Not just networks are vulnerable
- **Passive wiretapping** is listening in illicitly on conversations
- **Active wiretapping** is injecting traffic illicitly

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Wiretapping on Wires

- Signals can be trapped at many points
- Actually tapping into some physical wires is possible
- Other “wires” are broadcast media
 - **Packetsniffers** can listen to all traffic
- Subverted routers and gateways also offer access

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Wiretapping on Wireless

- Often just a matter of putting an antenna up
 - Though position may matter a lot
 - Generally not even detectable that it’s happening
- Active threats are easier to detect
 - And, for satellites, technically challenging

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Impersonation

- A packet comes in over the network
 - With some source indicated in its header
- Often, the action to be taken with the packet depends on the source
- But attackers may be able to create packets with false sources

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Methods of Network Impersonations

- Even in standard protocols, often easy to change fields in a header
 - When created or later
 - E.g., IP allows forging “from” addresses
- Existing networks have little or no built-in authentication

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Authentication to Foil Impersonation

- Higher level protocols often require authentication of transmissions
- Much care required to ensure proper authentication
- And not having authentication underneath can cause many problems
- Authentication schemes are rarely perfect

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Violations of Message Confidentiality

- Other problems can cause messages to be inappropriately divulged
- Misdelivery can send a message to the wrong place
 - Clever attackers can make it happen
- Message can be read at an intermediate gateway or a router
- Sometimes an intruder can get useful information just by traffic analysis

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Message Integrity

- Even if the attacker can't create the packets he wants, sometimes he can alter proper packets
- To change the effect of what they will do

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Methods of Attacks on Message Integrity

- Replacing part of a packet
- Changing headers to alter destination of a packet
 - Or its source
- Inserting improper packets into a proper packet stream

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Denial of Service

- Attacks that prevent legitimate users from doing their work
- By flooding the network
- Or corrupting routing tables
- Or flooding routers
- Or destroying key packets

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How Do Denial of Service Attacks Occur?

- Basically, the attacker injects some form of traffic
- Most current networks aren't built to throttle uncooperative parties very well
- All-inclusive nature of the Internet makes basic access trivial
- Universality of IP makes reaching most of the network easy

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Some Sample Attacks

- Smurf attacks
- SYN flood
- Ping of Death

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Smurf Attacks

- Attack on vulnerability in IP broadcasting
- Send a ping packet to IP broadcast address
 - With forged “from” header of your target
- Resulting in a flood of replies from the sources to the target
- Easy to fix at the intermediary
 - Don't allow IP broadcasts to originate outside your network
- No good solutions for victim

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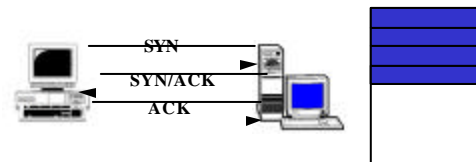
SYN Flood

- Based on vulnerability in TCP
- Attacker uses initial request/response to start TCP session to fill a table at the server
- Preventing new real TCP sessions
- SYN cookies and firewalls with massive tables are possible defenses

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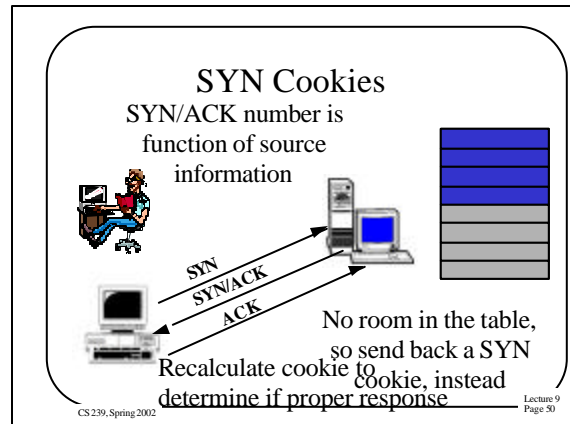
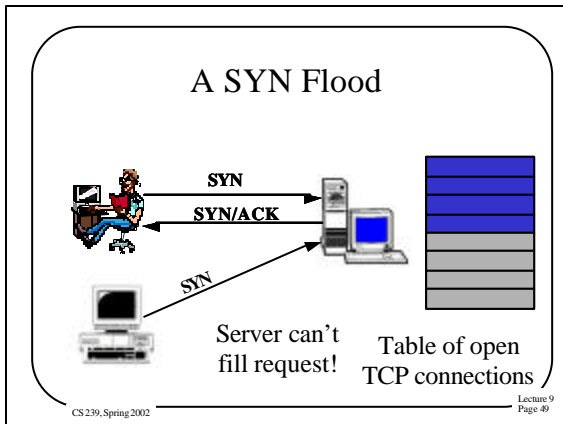
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Normal SYN Behavior



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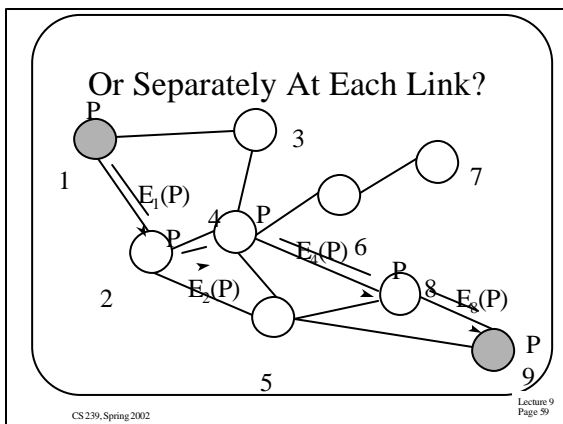
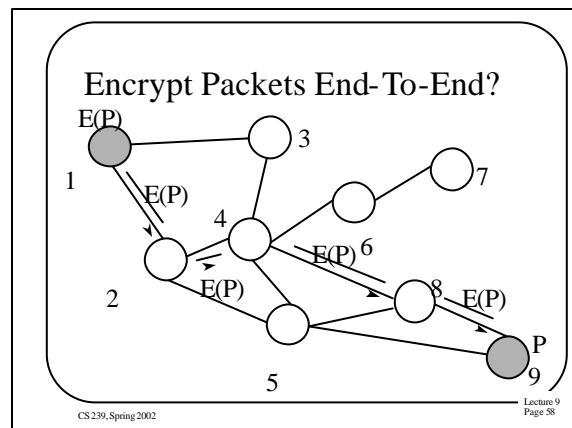
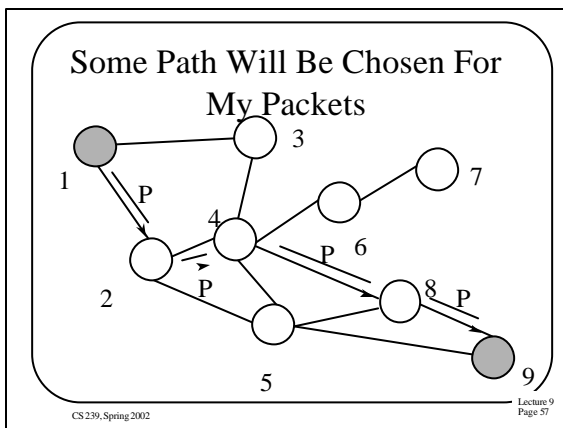
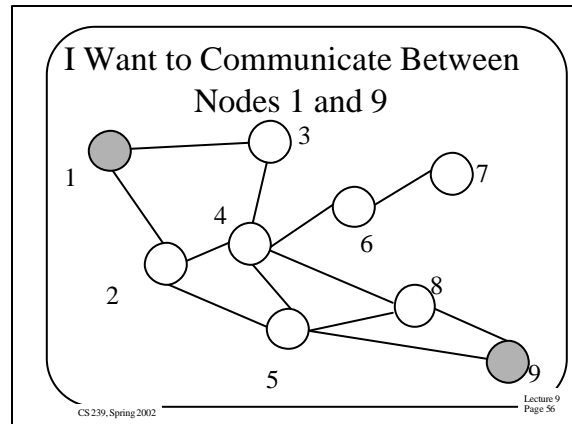
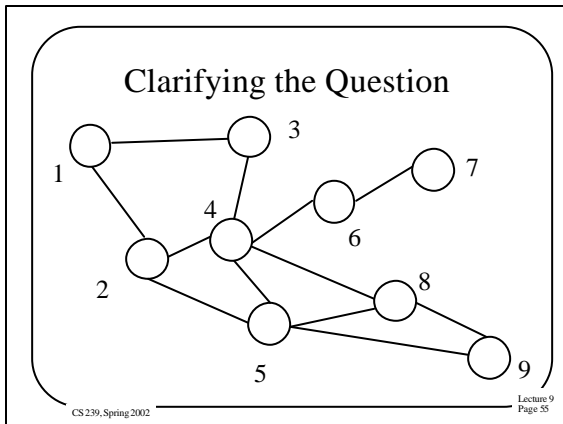


- ### The Ping of Death
- IP packets are supposed to be no longer than 65,535 bytes long
 - Can improperly send longer IP packets
 - Some OS networking software wasn't prepared for that
 - Resulting in buffer overflows and crashes
 - Can filter out pings, but other IP packets can also cause problem
 - OS patches really solve the problem
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- ### Network Security Mechanisms
- Again, the usual suspects -
 - Encryption
 - Authentication
 - Access control
 - Data integrity mechanisms
 - Traffic control
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- ### Encryption for Network Security
- Relies on the kinds of encryption algorithms and protocols discussed previously
 - But network security tends to only worry about the data transport issues
 - Which leads to an important question -
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- ### Link Encryption vs. End-to-End Encryption
- Should encryption be applied between pairs of hosts?
 - Or should encryption be applied between the endpoints of applications?
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- ### Well, What Difference Does It Make?
- The two methods have very different characteristics
 - Level of user/application involvement
 - Scaling properties
 - Trust requirements
 - Adaptability of transmission
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Link Level Encryption

- + Transparent to the user
- + Scaling related to number of links
- + Limits encryption to where it's needed
- + Can adapt data in transit
- Not as much user/application control
- May be applied unnecessarily
- Must trust intermediate nodes

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End-To-End Encryption

- + Greater possibilities for user control
- + Need not trust network components
- + Easier to apply selectively
- More user/application intervention required
- Data stream can't be adapted (much)
- Scaling related to logical connections

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Authentication for Network Security

- Various entities need to be authenticated
 - Hosts to hosts
 - Users to hosts
 - Hosts to users
- Because of inherent insecurities of networks, cryptographic methods used

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Access Control

- When a node is put on a network, potentially all its resources become available over the network
- How do we control who can access resources?
- And how?

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Data Integrity Mechanisms

- Bad things can happen if attackers can change data values
 - Either while in transit in the net
 - Or by remotely accessing a machine
- How do we keep our data intact?

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Checksums, Secure Hashes, and Digital Signatures

- Checksums can tell us if the data has changed
 - If the checksum hasn't been altered
- Secure hashes use cryptographic techniques
 - If the hash is protected
- Digital signatures provide full protection
 - At full cryptographic costs

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Traffic Control Mechanisms

- Filtering
 - Ingress filtering
 - Egress filtering
- Protection against traffic analysis
 - Padding
 - Routing control

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Ingress Filtering

- As packets enter router/switch/firewall, apply filtering rules
- Typically, drop packets not meeting some criteria
- Common example is firewall filtering
- Ingress filtering can help detect packets with bad “from” addresses
 - But only if forged address is “inside”

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Egress Filtering

- Routers/switches/firewalls filter packets leaving them
- To catch packets likely to cause trouble
- Egress filtering is commonly prescribed to handle forged “from” addresses
 - Only let out packets with “from” addresses in your domain
 - But not widely used
 - Since it provides few benefits to its user

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Padding

- Sometimes you don’t want intruders to know what your traffic characteristics are
- Padding adds extra traffic to hide the real stuff
- Requires that fake traffic is not differentiable from real
- Usually means encrypt it all
- Must be done carefully, or clever attackers can tell the good stuff from the noise

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Routing Control

- Use ability to control message routing to conceal the traffic in the network
- Especially important when trying to handle **covert channels**
 - Encapsulated users or programs trying to leak information out

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