

More on Network Security
CS 239
Security for Networks and
System Software
April 31, 2002

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Outline

- IPSEC
- Firewalls
- Virtual private networks

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IPSEC

- Until recently, the IP protocol had no standards for how to apply security
- Encryption and authentication layered on top
 - Or provided through ad hoc extensions
- Increasing security needs mandated a standard method of securing IP traffic

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How Was This Handled?

- The usual way that enhancements to standard Internet protocols are handled
 - The RFC/IETF mechanism
- Smart people worked out a proposal
- They published the proposal and requested comments
- Eventually agreement was reached

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IP-Security RFCs

- RFC 1825
 - Security Architecture for the Internet Protocol
- RFC 1826
 - IP Authentication Header
- RFC 1827
 - IP Encapsulating Security Payload

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Other Related RFCs

- RFC 1828 - IP Authentication Using Keyed MD5
- RFC 1829 - The ESP DES-CBC Transform
- RFC 1851 - The ESP Triple DES Transform
- RFC 1852 - IP Authentication Using Keyed SHA
- RFC 2085 - HMAC-MD5 IP Authentication With Replay Prevention
- And many, many others

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RFC 1825

- Defined the basics of security for the Internet Protocol
- Briefly, add per-packet encryption and authentication standards
- Basically, two mechanisms
 - A way to authenticate IP packets
 - A way to encrypt IP packets

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What Is Covered

- Message integrity
- Message authentication
- Message confidentiality

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What Isn't Covered

- Non-repudiation
- Digital signatures
- Key distribution
- Traffic analysis
- Handling of security associations
- Some of these covered in later RFCs and related standards

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Some Important Terms for IPSEC

- Security Association - “A set of security information related to a given network connection or set of connections”
 - Basically, a secure channel
- SPI (Security Parameters Index) - “An unstructured opaque index which is used in conjunction with the Destination Address to identify a particular Security Association”
 - Basically, a unique identifier

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General Structure of IPSEC

- Really designed for end-to-end encryption
 - Though could do link level
- Designed to operate with either IPv4 or IPv6
- Meant to operate with a variety of different encryption protocols
- And to be neutral to key distribution methods

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Security Associations

- Groups of entities that legitimately are cooperating in use of IPSEC for a particular connection
 - Hosts, applications, gateways, etc.
- Uniquely identified by:
 - Destination address
 - Plus a Security Parameter Index
 - Basically a pseudo-random number

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Creating Security Associations

- Setting them up properly is a major task in itself
- Not covered in basic IPSEC RFCs
 - But being heavily studied
- Normally one way
 - Two way traffic requires two Security Associations

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New IPSEC Headers

- The RFCs define two new types of headers for IP packets
 - The Authentication Header
 - The Encapsulating Security Payload

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IP Authentication Header

- Provides integrity and authentication
 - Not confidentiality
- Calculated using all fields in the IP datagram
 - Except those that change in transit
 - So both data and headers are authenticated

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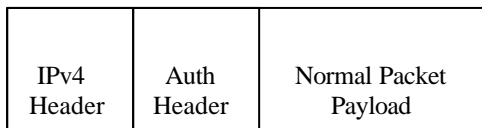
IP Authentication and Backwards Compatibility

- The authentication header is carried in the packet payload
- Non-participating routers can ignore it
- Participating routers know its payload location and can extract and check it as necessary

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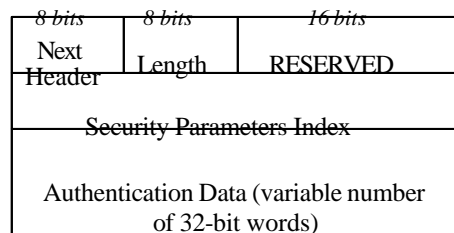
IPv4 Version of Packet With Authentication Header



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What's In the Authentication Header?



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Authentication Header Fields

- **Next header** identifies the next header in the packet
 - Possibly unrelated to IPSEC
- **Length** is length of this header's Authentication Data
- **Reserved** is, well, reserved
- **SPI** identifies the Security Association
- **Authentication data** is the actual "signature"

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IP Encapsulating Security Payload (ESP)

- Encrypt the data and place it within the ESP
- The ESP has normal IP headers
- Can be used to encrypt just the real data of the packet
- Or the entire IP packet

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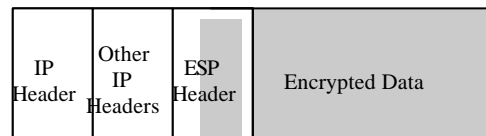
ESP Modes

- Tunnel mode
 - Original IP datagram is encrypted and placed in ESP
 - Unencrypted headers wrapped around ESP
- Transport mode
 - Encrypt just the transport-level data in the original packet
 - No IP headers encrypted

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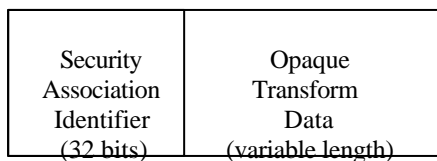
Secure IP Datagram Example



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The ESP Header



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Using ESP in Tunnel Mode

- Encrypt the IP datagram
 - The entire datagram
- Encapsulate it in a cleartext IP datagram
- Routers not understanding IPSEC can still handle it
- Receiver reverses the process

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ESP in Transport Mode

- Extract the transport-layer frame
 - E.g., TCP, UDP, etc.
- Encapsulate it in an ESP
- Encrypt it
- The encrypted data is now the last payload of a cleartext IP datagram

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What's the Status of IPSEC?

- The standard is done
- Widely implemented and used
 - In both Unix and Windows products
- The architecture doesn't require everyone to use it
- Generally considered to be a successful extension to IP

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What More Is Needed?

- Key distribution
 - E.g., IKE
- Security association handling
 - Also dealt with by IKE
- Implementations of IPSEC and IKE are freely available
- More work on broadcast/multicast use

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IPSEC and the AES Ciphers

- IPSEC is being adapted to use the new AES
- Currently, an Internet Draft memo describes using AES with IPSEC
- Further drafts looking at different modes/aspects of AES
- Expected that AES will become default for ESP in IPSEC

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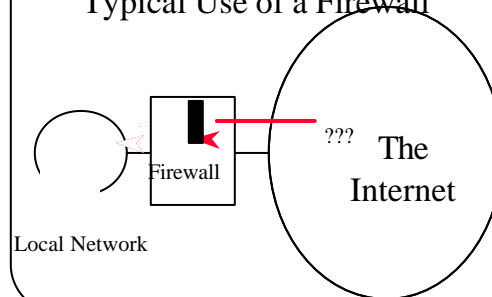
Firewalls

- “A system or combination of systems that enforces a boundary between two or more networks” - NCSA Firewall Functional Summary
- Usually, a computer that keeps the bad guys out

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Typical Use of a Firewall



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What Is a Firewall, Really?

- Typically a machine that sits between a LAN/WAN and the Internet
- Running special software
- That somehow regulates network traffic between the LAN/WAN and the Internet

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Firewalls and Perimeter Defense

- Firewalls implement a form of security called *perimeter defense*
- Protect the inside of something by defending the outside strongly
 - The firewall machine is often called a *bastion host*
- Control the entry and exit points
- If nothing bad can get in, I'm safe, right?

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Weaknesses of Perimeter Defense Models

- Breaching the perimeter compromises all security
- Windows passwords are a form of perimeter defense
 - If you get past the password, you can do anything
- Perimeter defense is part of the solution, not the entire solution

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Types of Firewalls

- Filtering gateways
 - AKA screening routers
- Circuit gateways
 - Also a kind of screening router
- Application level gateways
 - AKA proxy gateways
- Hybrid (complex) gateways

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Filtering Gateways

- Based on packet routing information
- Look at information in the incoming packets' headers
- Based on that information, either let the packet through or reject it

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Example Use of Filtering Gateways

- Allow particular external machines to telnet into specific internal machines
 - Denying telnet to other machines
- Or allow full access to some external machines
- And none to others

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A Fundamental Problem

- Today's IP packet headers aren't authenticated
 - And are pretty easy to forge
- If your filtering firewall trusts packet headers, it offers little protection
- Situation may be improved by IPSEC
 - But hasn't been yet

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One Exception to This Problem

- Checking internal addresses
- Safety procedures inside the security perimeter may limit some services to LAN members
- The firewall can check that incoming packets don't claim to be internal to the LAN

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Filtering Based on Ports

- Most incoming traffic is destined for a particular machine and port
 - Which can be derived from the IP and TCP headers
- Only let through packets to select machines at specific ports
- Makes it impossible to externally exploit flaws in little-used ports
 - If you configure the firewall right . . .

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Pros and Cons of Filtering Gateways

- + Fast
- + Cheap
- + Flexible
- + Transparent
- Limited capabilities
- Dependent on header authentication
- Generally poor logging
- May rely on router security

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Circuit Gateways

- Another kind of filtering firewall
- Used when internal machines request service from machines outside the firewall
- Makes it look like the request came from the firewall
 - Concealing internal system details

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Application Level Gateways

- Also known as proxy gateways
- Firewalls that understand the application-level details of network traffic
 - To some degree
- Traffic is accepted or rejected based on the probable results of accepting it

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How Application Level Gateways Work

- The firewall serves as a general framework
- Various proxies are plugged into the framework
- Incoming packets are examined
 - And handled by the appropriate proxy

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Firewall Proxies

- Programs capable of understanding particular kinds of traffic
 - E.g., FTP, HTTP, videoconferencing
- Proxies are specialized
- A good proxy must have deep understanding of the network application

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An Example Proxy

- A proxy to audit email
- What might such a proxy do?
 - Only allow email from particular hosts through
 - Or from particular users
 - Or filter out email with unsafe inclusions (like executables)

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What Are the Limits of Proxies?

- Proxies can only test for threats they understand
- Either they must permit a very limited set of operations
- Or they must have deep understanding of the program they protect
 - If too deep, they may share the flaw

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Pros and Cons of Application Level Gateways

- + Highly flexible
- + Good logging
- + Content-based filtering
- + Potentially transparent
- Slower
- More complex and expensive
- A good proxy is hard to find

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Hybrid Gateways

- A combination of two or more other types
 - Typically filtering gateways and proxy gateways
- Are they better?
 - If in parallel, no
 - If in series, maybe

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Firewall Characteristics

- Statefulness
- Transparency
- Firewalls and authentication
- Firewalls and encryption
- Firewalls and viruses

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Stateful Firewalls

- Much network traffic is connection-oriented
 - E.g., telnet and videoconferencing
- Proper handling of that traffic requires the firewall to maintain state
- But handling information about connections is more complex

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Firewalls and Transparency

- Ideally, the firewall should be invisible
 - Except when it vetoes access
- Users inside should be able to communicate outside without knowing about the firewall
- External users should be able to invoke internal services transparently

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Firewalls and Authentication

- Many systems want to allow specific sites or users special privileges
- Firewalls can only support that to the extent that strong authentication is available
 - At the granularity required
- For general use, may not be possible
 - In current systems

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Firewalls and Encryption

- Firewalls provide no confidentiality
 - For data they pass back and forth
- Unless the data is encrypted
- But if the data is encrypted, the firewall can't examine it
- So typically the firewall must be able to decrypt
 - Or only work on unencrypted parts of packets

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Firewalls and Link Encryption

- Inter-firewall encryption is essentially link level encryption
 - With all inherent problems
 - Except (presumably) that only trusted machines encrypt and decrypt
- More encryption can be applied at the application level
 - Limiting the firewall's options

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Firewalls and Viruses

- Firewalls are an excellent place to check for viruses
- Virus detection software can be run on incoming executables
- Requires that firewall knows when executables come in
- And must be reasonably fast

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Firewall Configuration and Administration

- Again, the firewall is the point of attack for intruders
- Thus, it must be extraordinarily secure
- How do you achieve that level of security?

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Firewall Hardening

- Devote a special machine only to firewall duties
- Alter OS operations on that machine
 - To allow only firewall activities
 - And to close known vulnerabilities
- Strictly limit access to the machine
 - Both login and remote execution

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Firewalls and Logging

- The firewall is the point of attack for intruders
- Logging activities there is thus vital
- The more logging, the better
- Should log what the firewall allows
- And what it denies
- Tricky to avoid information overload

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Closing the Back Doors

- Firewall security is based on assumption that all traffic goes through the firewall
- So be careful with:
 - Modem connections
 - Wireless connections
 - Portable computers
- Put a firewall at every entry point to your network
- And make sure all your firewalls are up to date

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Virtual Private Networks

- VPNs
- What if your company has more than one office?
- And they're far apart?
 - Like on opposite coasts of the US
- How can you have secure cooperation between them?

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Leased Line Solutions

- Lease private lines from some telephone company
- The phone company ensures that your lines cannot be tapped
 - To the extent you trust in phone company security
- Can be expensive and limiting

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Another Solution

- Communicate via the Internet
 - Getting full connectivity, bandwidth, reliability, etc.
 - At a lower price, too
- But how do you keep the traffic secure?
- Encrypt everything!

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Encryption and Virtual Private Networks

- Use encryption to convert a shared line to a private line
- Set up a firewall at each installation's network
- Set up shared encryption keys between the firewalls
- Encrypt all traffic using those keys

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Is This Solution Feasible?

- A VPN can be half the cost of leased lines (or less)
- And give the owner more direct control over the line's security
- If IPSEC succeeds, deployment and interoperation should be easy

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Key Management and VPNs

- All security of the VPN relies on key secrecy
- How do you communicate the key?
 - In early implementations, manually
 - Modern VPNs use something like IKE
- How often do you change the key?
 - IKE allows frequent changes

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VPNs and Firewalls

- VPN encryption is typically done between firewall machines
- Do I need the firewall for anything else?
- Probably, since I still need to allow non-VPN traffic in and out

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