

Security Protocols
CS 236
On-Line MS Program
Networks and Systems Security
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Outline

- Designing secure protocols
- Key exchange protocols
- Common security problems in protocols

Basics of Security Protocols

- Assume (usually) that your encryption is sufficiently strong
- Given that, how do you design a message exchange to achieve a given result securely?
- Not nearly as easy as you probably think
- Many of the concepts are important in many areas of computer/network security

Security Protocols

- A series of steps involving two or more parties designed to accomplish a task with suitable security
- Sequence is important
- Cryptographic protocols use cryptography
- Different protocols assume different levels of trust between participants

Types of Security Protocols

- Arbitrated protocols
 - Involving a trusted third party
- Adjudicated protocols
 - Trusted third party, after the fact
- Self-enforcing protocols
 - No trusted third party

Participants in Security Protocols



Alice



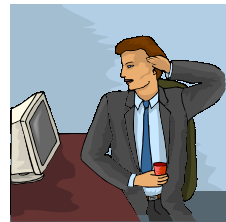
Bob

And the Bad Guys



Eve

Who only listens
passively



And sometimes
Alice or Bob
might cheat



Mallory

Who is actively
malicious

Trusted Arbitrator



Trent

A disinterested third party trusted by all legitimate participants

Arbitrators often simplify protocols, but add overhead and may limit applicability

Goals of Security Protocols

- Each protocol is intended to achieve some very particular goal
 - Like setting up a key between two parties
- Protocols may only be suitable for that particular purpose
- Important secondary goal is minimalism
 - Fewest possible messages
 - Least possible data
 - Least possible encryption

Key Exchange Protocols

- Often we want a different encryption key for each communication session
- How do we get those keys to the participants?
 - Securely
 - Quickly
 - Even if they've never communicated before

Key Exchange With Symmetric Encryption and an Arbitrator

- Alice and Bob want to talk securely with a new key
- They both trust Trent
 - Assume Alice & Bob each share a key with Trent
- How do Alice and Bob get a shared key?

Step One



Alice

K_A



Bob

K_B

*Alice
Requests
Session
Key for
Bob*



Trent

K_A

K_B

Who knows
what at this
point?

Step Two



K_A

Alice

$E_{K_A}(K_S),$
 $E_{K_B}(K_S)$



K_B

Bob

Who knows
what at this
point?



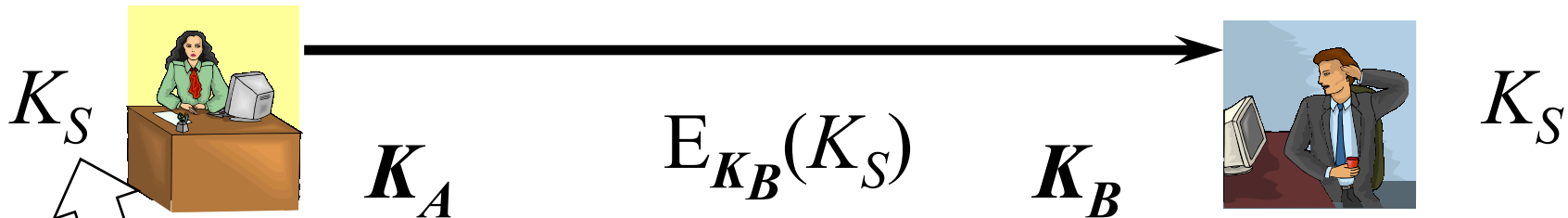
$E_{K_A}(K_S),$
 $E_{K_B}(K_S)$

K_A

Trent
 K_S

K_B

Step Three



Alice
 $E_{K_A}(K_S),$
 $E_{K_B}(K_S)$

Bob

Who knows
what at this
point?



K_A Trent K_B
 K_S

What Has the Protocol Achieved?

- Alice and Bob both have a new session key
- The session key was transmitted using keys known only to Alice and Bob
- Both Alice and Bob know that Trent participated
- But there are vulnerabilities

Problems With the Protocol

- What if the initial request was grabbed by Mallory?
- Could he do something bad that ends up causing us problems?
- Yes!

The Man-in-the-Middle Attack

- A class of attacks where an active attacker interposes himself secretly in a protocol
- Allowing alteration of the effects of the protocol
- Without necessarily attacking the encryption

Applying the Man-in-the-Middle Attack



Alice

*Alice
Requests
Session
Key for
Mallory*

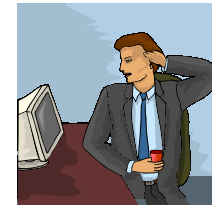
K_A



Mallory



Trent



Bob

More precisely,
who knows,
what do they think
what at this
they know?
point?

K_B

K_M

K_A

K_M

K_B

Trent Does His Job



Alice

K_A



K_M



K_B

Bob

Mallory

$E_{K_A}(K_S),$
 $E_{K_M}(K_S)$



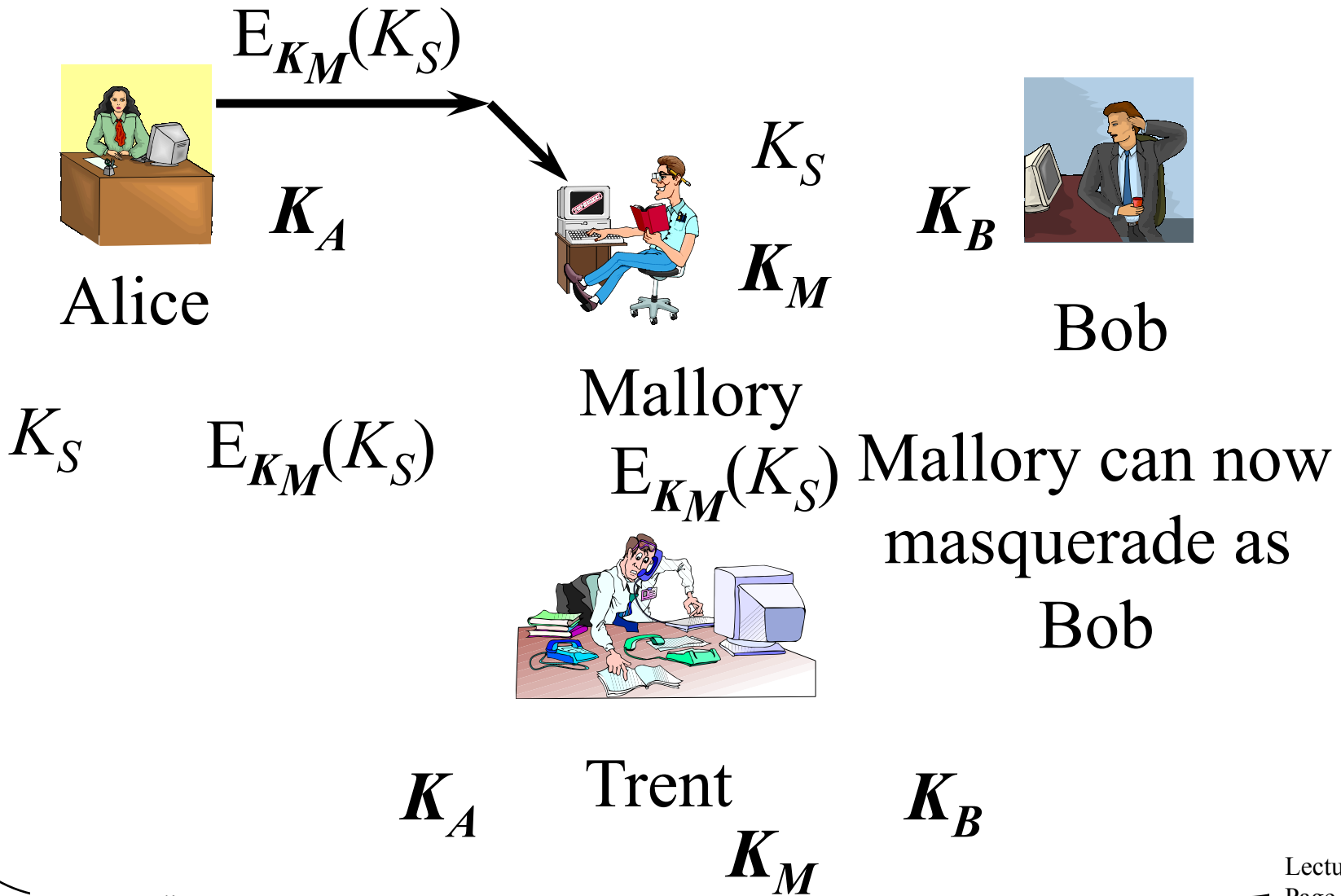
Trent

K_A

K_M

K_B

Alice Gets Ready to Talk to Bob



Really Getting in the Middle



Alice

K_A



K_M



Bob

K_{S1}

K_S

$E_{K_M}(K_{S1})$
 $E_{K_B}(K_{S1})$

Mallory

K_S
 K_{S1}

$E_{K_B}(K_{S1})$

Mallory can also
 ask Trent for a
 key to talk to
 Bob



Trent

K_A

K_M

K_B

Mallory Plays Man-in-the-Middle



Alice



Mallory



Bob

K_{S1}

K_S

Alice's big secret

$E_{K_S}(\text{Alice's big secret})$

$E_{K_S}(\text{Bob's big secret})$

Bob's big secret

K_S
 K_{S1}

$E_{K_S}(\text{Alice's big secret})$

$E_{K_{S1}}(\text{Alice's big secret})$

$E_{K_{S1}}(\text{Bob's big secret})$

$E_{K_S}(\text{Bob's big secret})$

Alice's big secret

Bob's big secret

Alice's big secret

Bob's big secret

$E_{K_{S1}}(\text{Bob's big secret})$

$E_{K_S}(\text{Bob's big secret})$

Defeating the Man In the Middle

- Problems:
 - 1). Trent doesn't really know what he's supposed to do
 - 2). Alice doesn't verify he did the right thing
- Minor changes can fix that
 - 1). Encrypt request with K_A
 - 2). Include identity of other participant in response - $E_{K_A}(K_S, \text{Bob})$

Applying the First Fix



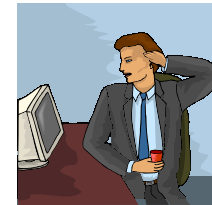
Alice

K_A



Mallory

K_M



Bob

K_B

E_{K_A} (Alice
Requests
Session
Key for
Bob)

Mallory can't
read the request

And Mallory
can't forge or
alter Alice's
request



Trent

K_A

K_B

K_M

But There's Another Problem

- A replay attack
- Replay attacks occur when Mallory copies down a bunch of protocol messages
- And then plays them again
- In some cases, this can wreak havoc
- Why does it here?

Step One



Alice

K_A



Bob

K_B



Mallory

E_{K_A} (Alice Requests Session Key for Bob)

E_{K_A} (Alice Requests Session Key for Bob)



Trent

K_A

K_B

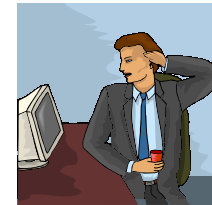
Step Two



Alice

K_A

$E_{K_A}(K_S),$
 $E_{K_B}(K_S)$



Bob

K_B



Mallory

E_{K_A} (Alice
Requests
Session
Key for
Bob)

$E_{K_A}(K_S),$
 $E_{K_B}(K_S)$



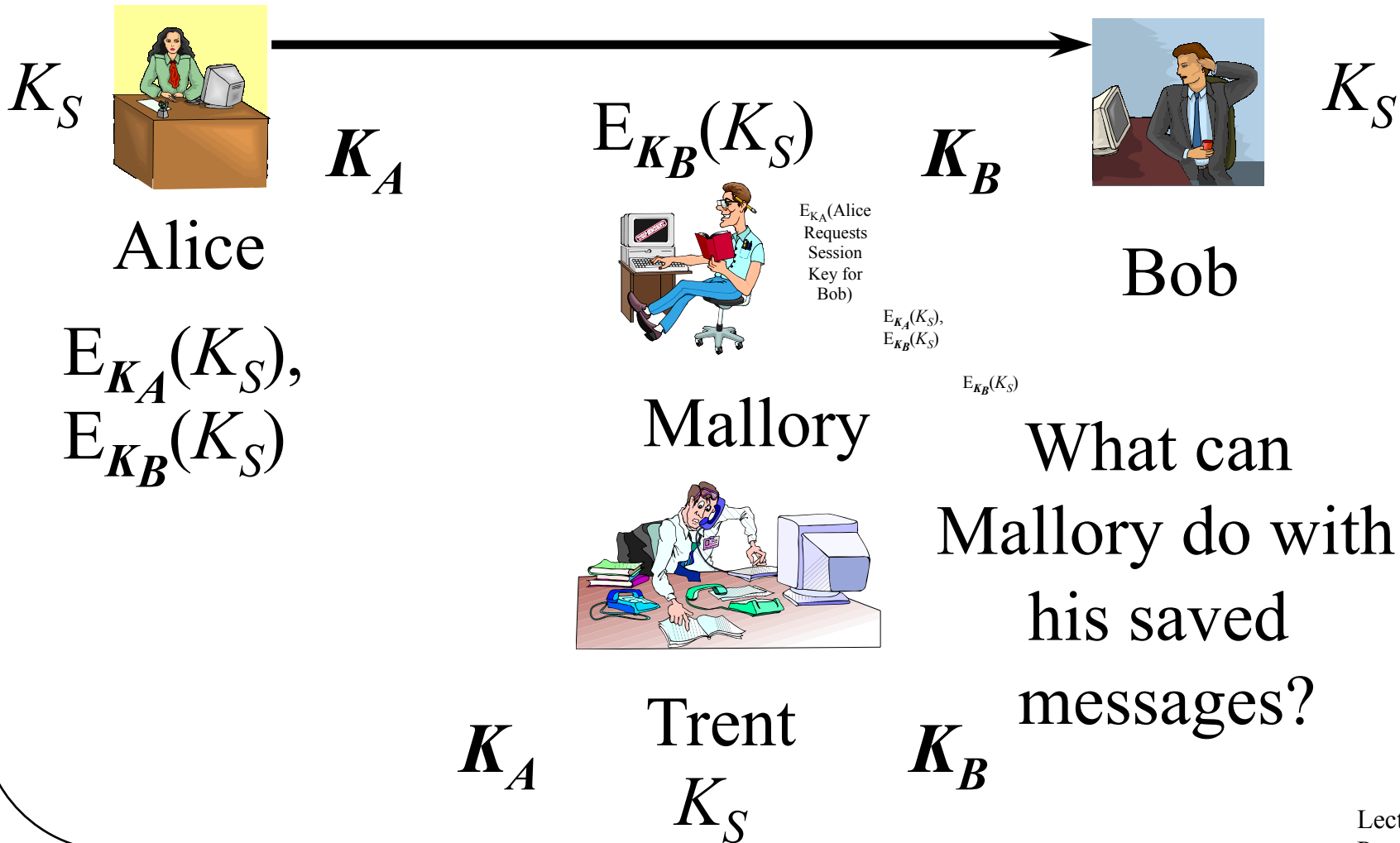
Trent

K_A

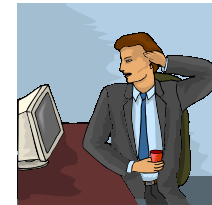
K_S

K_B

Step Three



Mallory Waits for His Opportunity



E_{K_A} (Alice Requests Session Key for Bob)

K_A

K_B

$E_{K_A}(K_S),$
 $E_{K_B}(K_S)$

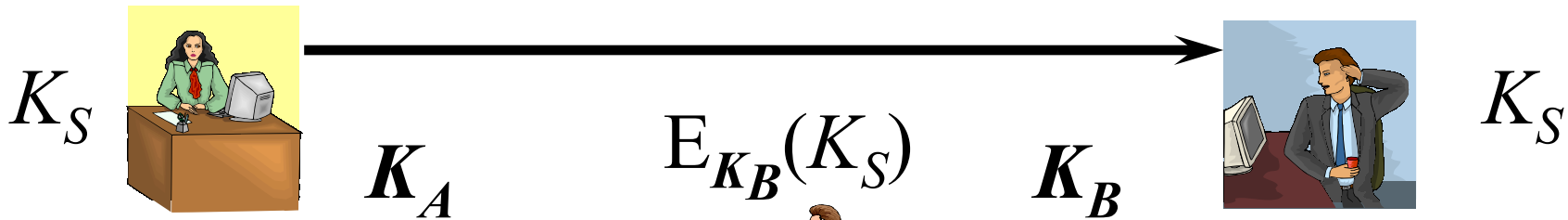
E_{K_A} (Alice Requests Session Key for Bob)

$E_{K_A}(K_S),$
 $E_{K_B}(K_S)$

$E_{K_B}(K_S)$

Mallory

What Will Happen Next?



E_{K_A} (Alice Requests Session Key for Bob)

$E_{K_A}(K_S), E_{K_B}(K_S)$

$E_{K_B}(K_S)$

K_S

Mallory



What's so bad about that?

What if Mallory has cracked K_S ?

K_A

K_B