

DNS Security

- The Domain Name Service (DNS) translates human-readable names to IP addresses
 - E.g., `thesiger.cs.ucla.edu` translates to `131.179.192.144`
 - DNS also provides other similar services
- It wasn't designed with security in mind

DNS Threats

- Threats to name lookup secrecy
 - Definition of DNS system says this data isn't secret
- Threats to DNS information integrity
 - Very important, since everything trusts that this translation is correct
- Threats to DNS availability
 - Potential to disrupt Internet service

What Could Really Go Wrong?

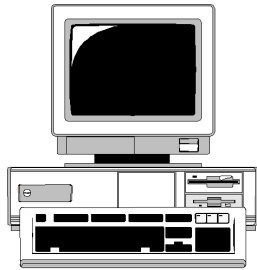
- DNS lookups could be faked
 - Meaning packets go to the wrong place
- The DNS service could be subject to a DoS attack
 - Or could be used to amplify one
- Attackers could “bug” a DNS server to learn what users are looking up

Where Does the Threat Occur?

- Unlike routing, threat can occur in several places
 - At DNS servers
 - But also at DNS clients
 - Which is almost everyone
- Core problem is that DNS responses aren't authenticated

The DNS Lookup Process

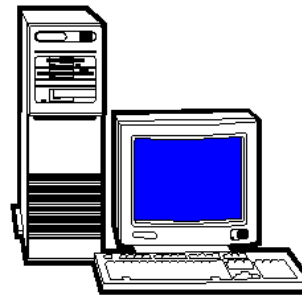
```
lookup thesiger.cs.ucla.edu
```



```
ping thesiger.cs.ucla.edu
```

Should result in a ping
packet being sent to
131.179.191.144

```
answer 131.179.191.144
```



If the answer is
wrong, in standard
DNS the client is
screwed

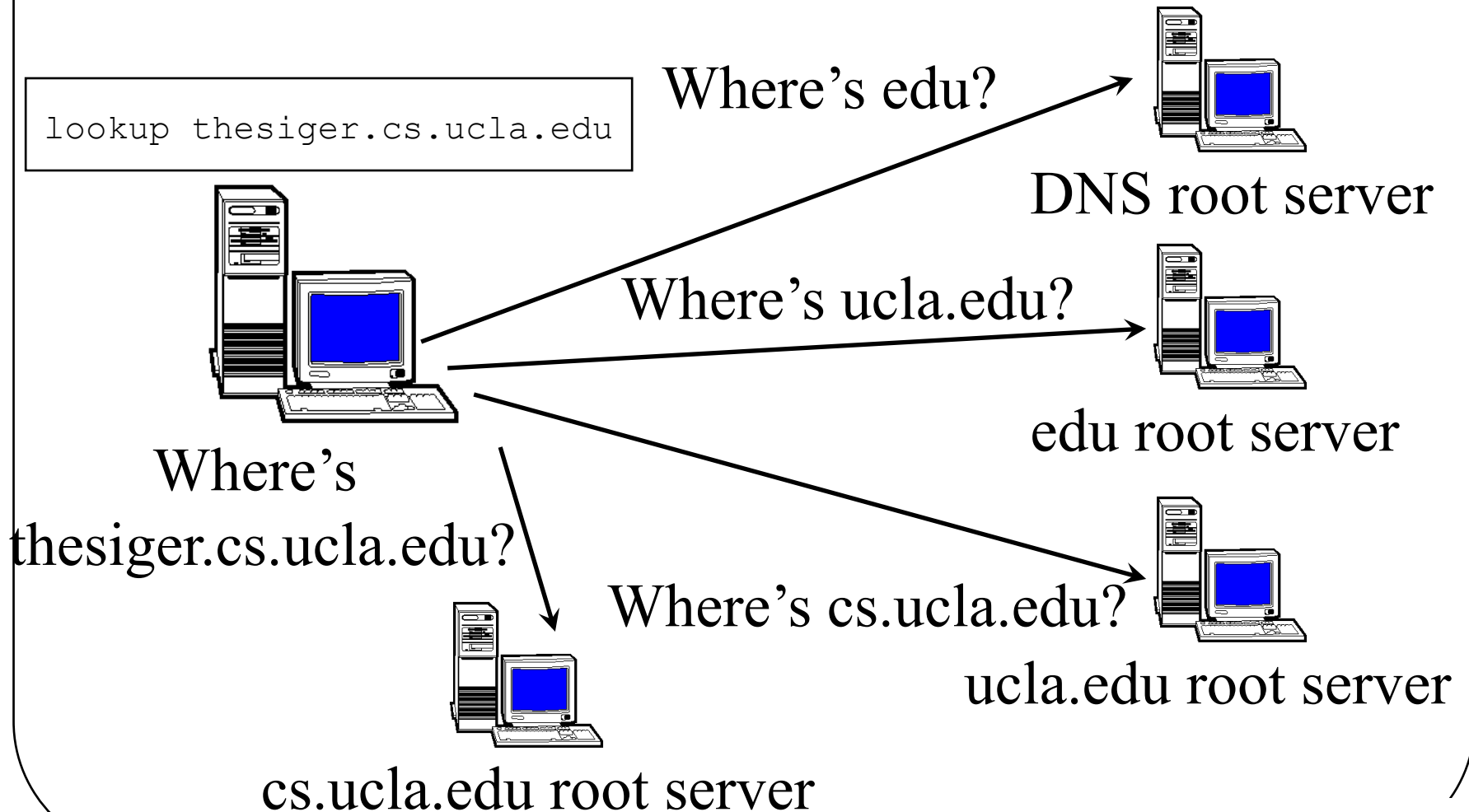
How Did the DNS Server Perform the Lookup?

- Leaving aside details, it has a table of translations between names and addresses
- It looked up `thesiger.cs.ucla.edu` in the table
- And replied with whatever the address was

Where Did That Table Come From?

- Ultimately, the table entries are created by those owning the domains
 - On a good day . . .
- And stored at servers that are authoritative for that domain
- In this case, the UCLA Computer Science Department DNS server ultimately stored it
- Other servers use a hierarchical lookup method to find the translation when needed

Doing Hierarchical Translation



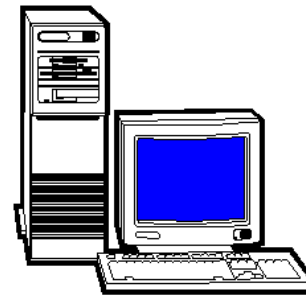
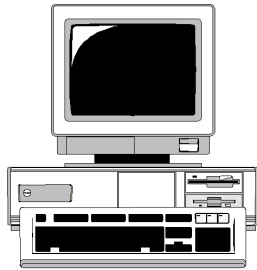
Where Can This Go Wrong?

- Someone can spoof the answer from a DNS server
 - Relatively easy, since UDP is used
- One of the DNS servers can lie
- Someone can corrupt the database of one of the DNS servers

The Spoofing Problem

```
lookup thesiger.cs.ucla.edu
```

```
answer 131.179.191.144
```



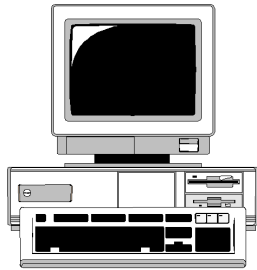
Unfortunately,
most DNS stub
resolvers will
take the first
answer

```
answer 97.22.101.53
```

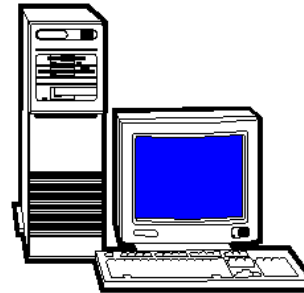


DNS Servers Lying

```
lookup thesiger.cs.ucla.edu
```



```
answer 97.22.101.53
```

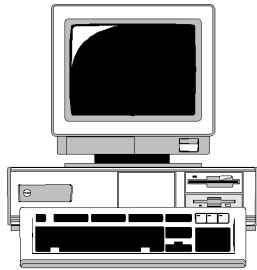


...	...
...	...
...	...
...	...
...	...
thesiger.cs.ucla.edu	131.178.192.144
...	...
...	...
...	...
...	...
...	...

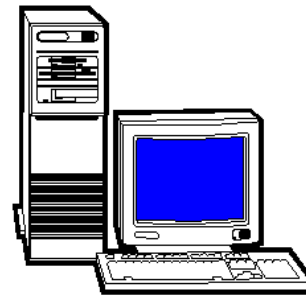
That wasn't very nice of him!

DNS Database Corruption

```
lookup thesiger.cs.ucla.edu
```



```
answer 97.22.101.53
```



...	...
...	...
...	...
...	...
...	...
thesiger.cs.ucla.edu	97.22.101.53
...	...
...	...
...	...
...	...
...	...

The DNSSEC Solution

- Sign the translations
- Who does the signing?
 - The server doing the response?
 - Or the server that “owns” the namespace in question?
- DNSSEC uses the latter solution

Implications of the DNSSEC Solution

- DNS databases must store signatures of resource records
- There must be a way of checking the signatures
- The protocol must allow signatures to be returned

Checking the Signature

- Basically, use certificates to validate public keys for namespaces
- Who signs the certificates?
 - The entity controlling the higher level namespace
- This implies a hierarchical solution

The DNSSEC Signing Hierarchy

- In principle, ICANN signs for itself and for top level domains (TLDs)
 - Like .com, .edu, country codes, etc.
- Each TLD signs for domains under it
- Those domains sign for domains below them
- And so on down

An Example

- Who signs the translation for `thesiger.cs.ucla.edu` to `131.179.192.144`?
- The UCLA CS DNS server
- How does someone know that's the right server to sign?
- Because the UCLA server says so
 - Securely, with signatures
- The edu server verifies the UCLA server's signature
- Ultimately, hierarchical signatures leading up to ICANN's attestation of who controls the edu namespace
- Where do you keep that information?
 - In DNS databases

Using DNSSEC

- To be really secure, you must check signatures yourself
- Next best is to have a really trusted authority check the signatures
 - And to have secure, authenticated communications between trusted authority and you

A Major Issue

- When you look up something like `cs.ucla.edu`, you get back a signed record
- What if you look up a name that doesn't exist?
- How can you get a signed record for every possible non-existent name?

The DNSSEC Solution

- Names are alphabetically orderable
- Between any two names that exist, there are a bunch of names that don't
- Sign the whole range of non-existent names
- If someone looks one up, give them the range signature

For Example,

•
•
•



lasr.cs.ucla.edu	131.179.192.136	
pelican.cs.ucla.edu	131.179.128.17	
pelican.uslacedu.edu	NOT ASSIGNED	
pelican.ucla.edu	131.179.128.16	
pelican.cs.ucla.edu	131.179.128.17	
toucan.cs.ucla.edu	131.179.128.16	

•
•
•

You get authoritative information that the name isn't assigned

Foils spoofing attacks

```
> host last.cs.ucla.edu
```

Status of DNSSEC

- Working implementations available
- In use in some places
- Heavily promoted
 - First by DARPA
 - Now by DHS
- Beginning to get out there

Status of DNSSEC Deployment

- ICANN has signed the root
 - .com, .gov, .edu, .org, .net all signed at top level
 - Not everyone below has signed, though
- Many “islands” of DNSSEC signatures
 - Signing for themselves and those below them
 - In most cases, just for themselves
- Utility depends on end machines checking signatures

Using DNSSEC

- Actually installing and using DNSSEC not quite as easy as it sounds
- Lots of complexities down in the weeds
- Particularly hard for domains with lots of churn in their namespace
 - Every new name requires big changes to what gets signed

Other DNS Solutions

- Encrypt communications with DNS servers
 - Prevents DNS cache poisoning
 - But assumes that DNS server already has right record
- Ask multiple servers
 - Majority rules or require consensus

Conclusion

- Correct Internet behavior depends on a few key technologies
 - Especially routing and DNS
- Initial (still popular) implementations of those technologies are not secure
- Work is ongoing on improving their security