

Operating System Security, Continued CS 239 Computer Security February 28, 2007

CS 236, Winter 2007

Lecture 11
Page 1

Outline

- Buffer overflows
- Designing secure operating systems
- Assuring OS security
- TPM and trusted computing
- Logging and auditing

CS 236, Winter 2007

Lecture 11
Page 2

Buffer Overflows

- One of the most common causes for compromises of operating systems
- Due to a flaw in how operating systems handle process inputs
 - Or a flaw in programming languages
 - Or a flaw in programmer training
 - Depending on how you look at it

CS 236, Winter 2007

Lecture 11
Page 3

What Is a Buffer Overflow?

- A program requests input from a user
- It allocates a temporary buffer to hold the input data
- It then reads all the data the user provides into the buffer, but . . .
- It doesn't check how much was provided

CS 236, Winter 2007

Lecture 11
Page 4

For Example,

```
int main(){
    char name[32];
    printf("Please type your name: ");
    gets(name);
    printf("Hello, %s", name);
    return (0);
}
```

- What if the user enters more than 32 characters?

CS 236, Winter 2007

Lecture 11
Page 5

Well, What If the User Does?

- The code continues reading data into memory
 - That's how `gets()` works
- The first 32 bytes go into `name`
- Where do the remaining bytes go?
- Onto the stack

CS 236, Winter 2007

Lecture 11
Page 6

Munging the Stack

- The temporary variable name is allocated on the stack
 - Close to the record of the function currently being run
- The overflow will spill into whatever's next on the stack
- Commonly, that's effectively going to overwrite the instruction pointer
 - Or the instruction pointer will be nearby

CS 236, Winter 2007

Lecture 11
Page 7

Using Buffer Overflows to Compromise Security

- Carefully choose what gets written into the instruction pointer
- So that the program jumps to something you want to do
 - Under the identity of the program that's running
- Such as, execute a command shell

CS 236, Winter 2007

Lecture 11
Page 8

Effects of Buffer Overflows

- Remote or unprivileged local user gets to run a program with greater privileges
- If buffer overflow is in a root program, gets all privileges, essentially
- Common mechanism to allow attackers to break into machines

CS 236, Winter 2007

Lecture 11
Page 9

Are Buffer Overflows Common?

- You bet!
- Weekly occurrences in major systems/applications
- Probably one of the most common security bugs

CS 236, Winter 2007

Lecture 11
Page 10

Some Recent Buffer Overflows

- Windows Media Player Plug-In
- Microsoft Windows Web Client
- LibPNG Graphics Library
- Metamail message processing
- Blackberry Enterprise Server
- And two others, just in last week's SANS vulnerability report

CS 236, Winter 2007

Lecture 11
Page 11

Fixing Buffer Overflows

- Check the length of the input
- Use programming languages that prevent them
- Put in OS controls that prevent overwriting the stack
- Put things in different places on the stack, making it hard to find the return pointer
- Why aren't these things commonly done?
- Presumably because programmers and designers neither know nor care about security

CS 236, Winter 2007

Lecture 11
Page 12

Desired Security Features of a Normal OS

- Authentication of users
- Memory protection
- File and I/O access control
- General object access control
- Enforcement of sharing
- Fairness guarantees
- Secure IPC and synchronization
- Security of OS protection mechanisms

CS 236, Winter 2007

Lecture 11
Page 13

Extra Features for a Trusted OS

- Mandatory and discretionary access control
- Object reuse protection
- Complete mediation
- Audit capabilities
- Intruder detection capabilities

CS 236, Winter 2007

Lecture 11
Page 14

How To Achieve OS Security

- Kernelized design
- Separation and isolation mechanisms
- Virtualization
- Layered design

CS 236, Winter 2007

Lecture 11
Page 15

Advantages of Kernelization

- Smaller amount of trusted code
- Easier to check every access
- Separation from other complex pieces of the system
- Easier to maintain and modify security features

CS 236, Winter 2007

Lecture 11
Page 16

Reference Monitors

- An important security concept for OS design
- A *reference monitor* is a subsystem that controls access to objects
 - It provides (potentially) complete mediation
- Very important to get this part right

CS 236, Winter 2007

Lecture 11
Page 17

Assurance of Trusted Operating Systems

- How do I know that I should trust someone's operating system?
- What methods can I use to achieve the level of trust I require?

CS 236, Winter 2007

Lecture 11
Page 18

Assurance Methods

- Testing
- Formal verification
- Validation

CS 236, Winter 2007

Lecture 11
Page 19

Secure Operating System Standards

- If I want to buy a secure operating system, how do I compare options?
- Use established standards for OS security
- Several standards exist

CS 236, Winter 2007

Lecture 11
Page 20

Some Security Standards

- U.S. Orange Book
- European ITSEC
- U.S. Combined Federal Criteria
- Common Criteria for Information Technology Security Evaluation

CS 236, Winter 2007

Lecture 11
Page 21

The U.S. Orange Book

- The earliest evaluation standard for trusted operating systems
- Defined by the Department of Defense in the late 1970s
- Now largely a historical artifact

CS 236, Winter 2007

Lecture 11
Page 22

Purpose of the Orange Book

- To set standards by which OS security could be evaluated
- Fairly strong definitions of what features and capabilities an OS had to have to achieve certain levels
- Allowing “head-to-head” evaluation of security of systems
 - And specification of requirements

CS 236, Winter 2007

Lecture 11
Page 23

Orange Book Security Divisions

- A, B, C, and D
 - In decreasing order of degree of security
- Important subdivisions within some of the divisions
- Requires formal certification from the government (NCSC)
 - Except for the D level

CS 236, Winter 2007

Lecture 11
Page 24

Some Important Orange Book Divisions and Subdivisions

- C2 - Controlled Access Protection
- B1 - Labeled Security Protection
- B2 - Structured Protection

CS 236, Winter 2007

Lecture 11
Page 25

The C2 Security Class

- Discretionary access control
 - At fairly low granularity
- Requires auditing of accesses
- And password authentication and protection of reused objects
- Windows NT has been certified to this class

CS 236, Winter 2007

Lecture 11
Page 26

The B1 Security Class

- Includes mandatory access control
 - Using Bell-La Padula model
 - Each subject and object is assigned a security level
- Requires both hierarchical and non-hierarchical access controls

CS 236, Winter 2007

Lecture 11
Page 27

The B3 Security Class

- Requires careful security design
 - With some level of verification
- And extensive testing
- Doesn't require formal verification
 - But does require "a convincing argument"
- Trusted Mach is in this class

CS 236, Winter 2007

Lecture 11
Page 28

The Common Criteria

- Modern international standards for computer systems security
- Covers more than just operating systems
- Design based on lessons learned from earlier security standards
- Lengthy documents describe the Common Criteria

CS 236, Winter 2007

Lecture 11
Page 29

Basics of Common Criteria Approach

- Something of an alphabet soup –
- The CC documents describe
 - The Evaluation Assurance Levels (EAL)
- The Common Evaluation Methodology (CEM) details guidelines for evaluating systems

CS 236, Winter 2007

Lecture 11
Page 30

Another Bowl of Common Criteria Alphabet Soup

- TOE – Target of Evaluation
- TSP – TOE Security Policy
 - Security policy of system being evaluated
- TSF – TOE Security Functions
 - HW, SW used to enforce TSP
- PP – Protection Profile
 - Implementation-dependent set of security requirements
- ST – Security Target
 - Predefined sets of security requirements

CS 236, Winter 2007

Lecture 11
Page 31

What's This All Mean?

- Highly detailed methodology for specifying :
 1. What security goals a system has
 2. What environment it operates in
 3. What mechanisms it uses to achieve its security goals
 4. Why anyone should believe it does so

CS 236, Winter 2007

Lecture 11
Page 32

TPM and Trusted Computing

- Can special hardware help improve OS security?
- Perhaps
- TPM is an approach to building such hardware
- The approach is commonly called “trusted computing”

CS 236, Winter 2007

Lecture 11
Page 33

What Is TPM?

- Special hardware built into personal computers
 - And other types of machines
- Tamperproof, special purpose
- Effective use requires interaction with software
 - Especially OS software
- Defined as a set of open standards

CS 236, Winter 2007

Lecture 11
Page 34

What Does TPM Hardware Do?

- Three basic core functionalities:
 - Secure storage and use of keys
 - Secure software attestations
 - Sealing data
- These functions can be used to build several useful security features

CS 236, Winter 2007

Lecture 11
Page 35

TPM Key Storage

- Keys are stored in a tamperproof area
- TPM hardware can generate RSA key pairs
 - Using true random number generator
- Each TPM chip has one permanent endorsement key
- Other keys generated as needed

CS 236, Winter 2007

Lecture 11
Page 36

The Endorsement Key

- Created when the chip was fabricated
- Used to sign attestations
 - To prove that this particular machine made the attestation
- A public/private key pair
 - Private part never leaves the trusted hardware

CS 236, Winter 2007

Lecture 11
Page 37

TPM Cryptography

- Some TPM hardware includes encryption and decryption functions
- To ensure keys are never outside a tamperproof perimeter

CS 236, Winter 2007

Lecture 11
Page 38

TPM Attestations

- Allows TPM to provide proof that a particular piece of software is running on the machine
 - An OS, a web browser, whatever
- Essentially, a signature on a hash of the software

CS 236, Winter 2007

Lecture 11
Page 39

An Example of an Attestation

- What version of Linux is running on this machine?
- TPM (with appropriate SW support) hashes the OS itself
- Signs the hash with its attestation key
- Sends the signature to whoever needs to know

CS 236, Winter 2007

Lecture 11
Page 40

Secure TPM Boot Facilities

- Use attestations to ensure that the boot loader is trusted code
- The trusted boot loader then checks the OS it intends to load
 - Trusted attestations can tell the boot loader if it's the right one
 - Bail out if it's not the right one
- Can prevent an attacker from getting you to boot a corrupted kernel

CS 236, Winter 2007

Lecture 11
Page 41

Sealing Data With TPM

- Encrypt the data with keys particular to one machine
 - Keys stored by TPM
- Data can only be decrypted successfully on that machine
- Can also seal storage such that only a particular application can access it

CS 236, Winter 2007

Lecture 11
Page 42

The TPM Controversy

- TPM can be used for many good security purposes
- But some believe it takes too much power from the user
 - E.g., can require user to prove he's running a particular browser before you give him a file
 - Or seal a file so only the owner's application can read it
- Many (but not all) critics worry especially about DRM uses
 - Also serious issues about companies using it to achieve anti-competitive effects
- Serious questions about practicality based on patching, various releases, etc.
 - Will you have to accept attestations for all of them?

CS 236, Winter 2007

Lecture 11
Page 43

Logging and Auditing

- An important part of a complete security solution
- Practical security depends on knowing what is happening in your system
- Logging and auditing is required for that purpose

CS 236, Winter 2007

Lecture 11
Page 44

Logging

- No security system will stop all attacks
 - Logging what has happened is vital to dealing with the holes
- Logging also tells you when someone is trying to break in
 - Perhaps giving you a chance to close possible holes

CS 236, Winter 2007

Lecture 11
Page 45

Access Logs

- One example of what might be logged for security purposes
- Listing of which users accessed which objects
 - And when and for how long
- Especially important to log failures

CS 236, Winter 2007

Lecture 11
Page 46

Other Typical Logging Actions

- Logging failed login attempts
 - Can help detect intrusions or password crackers
- Logging changes in program permissions
 - A common action by intruders
- Logging scans of ports known to be dangerous

CS 236, Winter 2007

Lecture 11
Page 47

Problems With Logging

- Dealing with large volumes of data
- Separating the wheat from the chaff
 - Unless the log is very short, auditing it can be laborious
- System overheads and costs

CS 236, Winter 2007

Lecture 11
Page 48

Log Security

- If you use logs to detect intruders, smart intruders will try to attack logs
 - Concealing their traces by erasing or modifying the log entries
- Append-only access control helps a lot here
- Or logging to hard copy
- Or logging to a remote machine

CS 236, Winter 2007

Lecture 11
Page 49

Local Logging vs. Remote Logging

- Should you log just on the machine where the event occurs?
- Or log it just at a central site?
- Or both?

CS 236, Winter 2007

Lecture 11
Page 50

Local Logging

- Only gives you the local picture
- More likely to be compromised by attacker
- Must share resources with everything else machine does
- Inherently distributed
 - Which has its good points and bad points

CS 236, Winter 2007

Lecture 11
Page 51

Remote Logging

- On centralized machine or through some hierarchical arrangement
- Can give combined view of what's happening in entire installation
- Machine storing logs can be specialized for that purpose
- But what if it's down or unreachable?
- A goldmine for an attacker, if he can break in

CS 236, Winter 2007

Lecture 11
Page 52

Desirable Characteristics of a Logging Machine

- Devoted to that purpose
 - Don't run anything else on it
- Highly secure
 - Control logins
 - Limit all other forms of access
- Reasonably well provisioned
 - Especially with disk

CS 236, Winter 2007

Lecture 11
Page 53

Auditing

- Security mechanisms are great
 - If you have proper policies to use them
- Security policies are great
 - If you follow them
- For practical systems, proper policies and consistent use are a major security problem

CS 236, Winter 2007

Lecture 11
Page 54

Auditing

- A formal (or semi-formal) process of verifying system security
- “You may not do what I expect, but you will do what I inspect.”
- A requirement if you really want your systems to run securely

CS 236, Winter 2007

Lecture 11
Page 55

Auditing Requirements

- Knowledge
 - Of the installation and general security issues
- Independence
- Trustworthiness
- Ideally, big organizations should have their own auditors

CS 236, Winter 2007

Lecture 11
Page 56

When Should You Audit?

- Periodically
- Shortly after making major system changes
 - Especially those with security implications
- When problems arise
 - Internally or externally

CS 236, Winter 2007

Lecture 11
Page 57

Auditing and Logs

- Logs are a major audit tool
- Some examination can be done automatically
- But part of the purpose is to detect things that automatic methods miss
 - So some logs should be audited by hand

CS 236, Winter 2007

Lecture 11
Page 58

A Typical Set of Audit Criteria

- For a Unix system
- Some sample criteria:
 - All accounts have passwords
 - Limited use of setuid root
 - Display last login date on login
 - Limited write access to system files
 - No “.” in PATH variables

CS 236, Winter 2007

Lecture 11
Page 59

What Does an Audit Cover?

- Conformance to policy
- Review of control structures
- Examination of audit trail (logs)
- User awareness of security
- Physical controls
- Software licensing and intellectual property issues

CS 236, Winter 2007

Lecture 11
Page 60

Does Auditing Really Occur?

- To some extent, yes
- 2005 CSI/FBI report says 87% of responding organizations did audits
 - Up from 82% in 2004
- Doesn't say much about the quality of the audits
- It's easy to do a bad audit